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Monitoring PC Hardware Sounds in Linux Systems Using the Daubechies D4 Wavelet.

Robert Karns Henry
East Tennessee State University

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Monitoring PC Hardware Sounds in Linux Systems Using the Daubechies D4 Wavelet

A thesis
presented to
the faculty of the Department of Computer Science
East Tennessee State University

In partial fulfillment
of the requirements for the degree
Master in Applied Computer Science

by
Robert K. Henry
December, 2005

David Tarnoff, Chair
Phillip Pfeiffer
Jeff Knisley

Keywords: Wavelet, Daubechies, DSP, Monitor, S.M.A.R.T., ALSA, Linux, failure prediction, remote, vibration analysis
ABSTRACT

Monitoring PC Hardware Sounds in Linux Systems Using the Daubechies D4 Wavelet

by

Robert K. Henry

Users of high availability (HA) computing require systems that run continuously, with little or no downtime. Modern PCs address HA needs by monitoring operating system parameters such as voltage, temperature, and hard drive status in order to anticipate possible system failure. However, one modality for PC monitoring that has been underutilized is sound. The application described here uses wavelet theory to analyze sounds produced by PC hard drives during standard operation. When twenty-nine hard drives were tested with the application and the results compared with the drives’ Self-Monitoring, Analysis, and Reporting Technology (S.M.A.R.T.) data, the binomial distribution’s low p-value of 0.012 indicated better than chance agreement. While the concurrence between the two systems shows that sound is an effective tool in detecting hardware failures, the disagreements between the systems show that the application can complement S.M.A.R.T. in an HA system.
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CHAPTER 1
INTRODUCTION

1.1 Motivation for Research

Users of high availability (HA) computing require systems that run continuously, with little or no downtime. One critical technology for achieving high availability is self-monitoring: the ability of a system to detect faults and signs of abnormalities. Self-monitoring is particularly important in environments such as massively parallel clusters, where downtime can be costly and disruptive, and the sheer mass of hardware makes the detection of faults and failures more complex and costly.

To make component failures easier to detect, modern computers often support data acquisition devices (DAQ) that monitor operating characteristics such as temperature, voltage, and fan speed. A more intelligent device, the Self-Monitoring, Analysis and Reporting Technology (S.M.A.R.T.) System, provides data on hard drive operation; basic input/output systems (BIOSes) and operating systems have used S.M.A.R.T. technology for years to identify abnormalities in hard drives.

One potentially useful but underutilized modality for detecting failing components in computers is sound. Sound has previously been used to detect abnormal system operation in automotive and mechanical engineering. Modern automobile engine controls, for instance, use an accelerometer to monitor vibrations caused by detonation, and then control spark advance accordingly.¹ Another project investigated using signals from accelerometers to monitor vibration in helicopter transmissions. In this project, accelerometers installed near critical

bearings and gears in the transmission transmit a signal that is analyzed for potential component failure.\(^2\)

Although most computer components are not mechanical, devices such as hard drives and cooling fans contain motors and bearings that may make abnormal sounds as they deteriorate. PC users often request service when they hear these sounds, which in turn are often used in troubleshooting.\(^3\) \(^4\)

In the case of cluster computing, the task of monitoring sound becomes much more difficult. There are three reasons for this. The tendency to house these systems in dedicated rooms precludes most users from hearing any sounds at all. Moreover, the number of PCs in these rooms makes it less likely that people who do come near the machines will hear abnormal noises, or if they do hear a sporadic strange noise, be able to identify the source of that noise.

\subsection*{1.2 Results}

The research described here produced a Linux-based prototype for a sound-based system for detecting and analyzing mechanical hard drive failures. This system uses a built-in sound card and a microphone to obtain digital samples of hardware sounds. Once acquired and digitized, sounds are decomposed using digital signal processing (DSP) techniques, and checked for symptoms of potential failure.

A design consideration was that the analysis system must present a minimal processing load on the system. This constraint was one consideration in choosing wavelet analysis as the DSP filter. Wavelets can provide useful filtering in time $O(n)$, where $n$ is the sample size: a significant performance improvement over some other DSP techniques.

\begin{thebibliography}{9}
\bibitem{3} Knowledge Base Online Help Topics, “Identifying and troubleshooting hard drive noise issues”, Maxtor Corporation, 2005, 
\bibitem{4} Service and Support, “How can I tell if the noise or sound my drive is making is normal?” , Western Digital Corporation, 2005, 
\end{thebibliography}
The prototype was tested on sounds from unserviceable hard drives. Digital recordings of the hard drives were played to the application and the results from the application captured for examination. The captured data consisted of $\chi^2$ comparisons between digitally filtered samples of the input sound and a template prepared to represent a normal quiescent hard drive. Those $\chi^2$ results exceeding a fixed threshold were used to classify hard drives as good or bad. The threshold was also studied to determine the optimum value for this parameter.

The results of the sound testing were compared to the S.M.A.R.T. data extracted from the hard drives. These two sets of observations were compared using Cohen’s kappa ($\kappa$), a statistical measure of observer agreement. Kappa corrects for agreement between the observers due to chance but says nothing about the validity of either set of observations. Though S.M.A.R.T. is being used to measure the validity of the sound measurement results, S.M.A.R.T. cannot be assumed as an infallible benchmark. This makes Kappa a useful comparison tool, because it measures the degree of agreement between two observers, either of which may be wrong. Comparing the S.M.A.R.T. observations with the sound observations with optimized threshold yielded a kappa as high as 0.3446, indicating a fair degree of agreement.

Agreement between the two observers was also compared using the binomial distribution. The binomial distribution is used to calculate the probability of the observers agreeing by chance. The p-value was computed for numbers of agreements equal or greater than observed numbers with a binomial probability of 50%. To test for significance, the p-value should be as low as possible. Though p-values below 0.05 are normally considered significant, the less restrictive value of 0.10 was used for this research because of the small number of drives available for test and because of inherent differences between sound and S.M.A.R.T. that would make perfect agreement impossible. The binomial distribution p-value testing for agreement between sound and S.M.A.R.T. by chance was as low as 0.012. This number indicates that there is a 0.012 or 1.2% probability that the agreement between the observers is by chance.

These results, in short, show that sound can be used as an effective indicator for hardware condition.
1.3 Overview of Thesis

The balance of this thesis is divided into five chapters. Chapter 2 discusses supporting technologies for this work, including data acquisition, signal processing, wavelets, sampling sound using the PC sound card under Linux, and methods of classifying wavelet data. Chapter 3 reviews the methodology and scope of the research. Chapter 4 discusses the design of the software components used in the system prototype. Chapter 5 examines the results of experiments with the prototype to assess the system’s effectiveness. Chapter 6 summarizes the thesis, offering suggestions for future work.
CHAPTER 2
BACKGROUND

This chapter discusses the technologies on which this research is based. Topics covered include hardware sensing technology for personal computers (PCs); the Self-Monitoring, Analysis and Reporting Technology (S.M.A.R.T.) system for analyzing hard drive performance; digital-signal-processing (DSP) techniques for sound analysis; and the Linux sound card technology used in this work for handling sound samples.

2.1 Hardware Sensors

2.1.1 Basic Sensors

Large computing systems have long incorporated sensors or data acquisition (DAQ) devices to monitor the operation of hardware. For example, in 1977, the Cray-1 computer had a monitoring system to detect power and cooling systems malfunctions.\(^5\)

In recent years, manufacturers have installed sensors in some personal computers (PCs) to monitor physical characteristics such as temperature, voltage, and other measurable quantities of the motherboard and peripheral devices. PC sensors obtain data on a system’s hardware and provide that data to the host system.

An example of a common, simple sensor is National Semiconductor’s LM75 temperature sensor.\(^6\) The LM75 incorporates a sensing element, an analog-to-digital converter (ADC), and other digital circuitry for reporting temperatures to the processor. The sensing element delivers a voltage signal to the ADC, which computes the digital number that represents the temperature. The host machine reads the digital temperature measurement from the device as a 9-bit 2’s-complement value representing temperatures from -55°C to 125°C with an accuracy of between ±2°C and ±3°C.


When installed in a computer, the LM75 is mapped to an I/O port that can be queried at any time by the host machine and the temperature read directly from the port. The LM75 can also be configured to generate an interrupt to the host computer if the measured temperature exceeds a specified value. In this configuration, attention from the host computer is required only when the system's temperature is of concern. In this case, interrupt handlers on the host computer must be designed to read the correct I/O port for the device and correctly interpret the data obtained. This means that the software must be designed for the specific computer or must be user configurable.

A more modern DAQ device, National Semiconductor’s multipurpose LM87 sensor, can acquire and digitize data from several different sources. The LM87 can measure temperatures from an on-chip sensor and accept input from two additional external temperature sensors. In addition, the device can sample voltages, measure and control the speeds of two sensor-equipped fans, and generate an interrupt when a parameter exceeds its tolerances.\(^7\) As in the case of the LM75, the LM87’s software controller must be specific to the host computer’s hardware.

2.1.2 Self-Monitoring, Analysis, and Reporting Technology (S.M.A.R.T.)

An additional type of PC sensor system is the Self-Monitoring, Analysis, and Reporting Technology (S.M.A.R.T.) found in ATA/IDE and SCSI hard drives.\(^8\) S.M.A.R.T. is based on IBM’s technology of Predictive Failure Analysis (PFA),\(^9\) a technology that IBM first provided to customers in 1992.\(^10\) S.M.A.R.T. sensors predict hard drive failure by measuring and analyzing a hard drive’s operating parameters including head flying height, drive spin-up time, and temperature. If head flying height is lower than normal, an above-average possibility of a head crash exists. A drive that takes longer than usual to spin up may have bearings that are wearing

---

\(^7\) Serial Interface System Hardware Monitor with Remote Diode Temperature Sensing, National Semiconductor Corporation, (9/11/2003), (LM87.pdf)

\(^8\) Mark Evans, SFF Committee Specification for Self-Monitoring, Analysis, and Reporting Technology, Quantum Corporation, (4/26/1996)


When the probability of failure exceeds a given threshold, the drive can warn of imminent failure, permitting backup of data and replacement of the component before failure occurs. Hughes, et. al. suggest that S.M.A.R.T. does not perform perfectly in reporting impending failures without excessive false alarms due to the low failure rate for hard drives, typically <1%/year.\textsuperscript{12}

Normally, a PC’s BIOS checks S.M.A.R.T. data during the power on self-test. Parameters that fall outside their threshold values cause the system to signal an error. In addition to the BIOS monitoring, several software applications are available for monitoring S.M.A.R.T. drive data. Most hard drive manufacturers provide software that reads, analyzes, or displays S.M.A.R.T. data for their own or other manufacturers’ drives, and make these applications available for download without cost from the manufacturer’s websites.

Vendor-supported and freeware applications are available that read, monitor, and display hard drive S.M.A.R.T. data. An example of a free S.M.A.R.T. application is smartmontools, released under the GNU General Public License and available for Windows as well as UNIX/Linux.\textsuperscript{13} Figure 1 presents a smartmontools report for a drive with an excessive raw read error rate count. This report indicates that the drive has failed and that its data should be backed up promptly.

The information presented in S.M.A.R.T. drive reports varies, depending on the drive manufacturer’s implementation of the S.M.A.R.T. standard. These variations may represent choices by manufacturers to ignore parameters that have not been shown to be useful in predicting failure for this particular device.

\textsuperscript{11} Playing it S.M.A.R.T.
\textsuperscript{13} Bruce Allen, smartmontools, \url{http://smartmontools.sourceforge.net} (2002-4)
smartctl version 5.33 [i386-pc-mingw32] Copyright (C) 2002-4 Bruce Allen
Home page is http://smartmontools.sourceforge.net/

--- START OF INFORMATION SECTION ---
Device Model: WDC AC22100H
Serial Number: WD-WM3610442859
Firmware Version: 10.07H1
User Capacity: 2,111,864,832 bytes
Device is: Not in smartctl database [for details use: -P showall]
ATA Version is: 1
ATA Standard is: Exact ATA specification draft version not indicated
Local Time is: Wed May 18 14:49:59 2005 Pacific Daylight Time
SMART is only available in ATA Version 3 Revision 3 or greater.
We will try to proceed in spite of this.
SMART support is: Ambiguous - ATA IDENTIFY DEVICE words 82-83 don't show if SMART supported.
Checking for SMART support by trying SMART ENABLE command.
SMART ENABLE appeared to work! Continuing.
SMART support is: Ambiguous - ATA IDENTIFY DEVICE words 85-87 don't show if SMART is enabled.
Checking to be sure by trying SMART RETURN STATUS command.
SMART support is: Enabled

--- START OF READ SMART DATA SECTION ---
SMART overall-health self-assessment test result: FAILED!
Drive failure expected in less than 24 hours. SAVE ALL DATA.
See vendor-specific Attribute list for failed Attributes.

General SMART Values:
Offline data collection status: (0x00) Offline data collection activity
was never started.
Auto Offline Data Collection: Disabled.
Total time to complete Offline
data collection: (900) seconds.
Offline data collection
capabilities: (0x03) SMART execute Offline immediate.
Auto Offline data collection on/off supp ort.
Suspend Offline collection upon new
command.
No Offline surface scan supported.
No Self-test supported.
No Conveyance Self-test supported.
No Selective Self-test supported.
SMART capabilities:
(0x0002) Does not save SMART data before
entering power-saving mode.
Supports SMART auto save timer.
Error logging capability:
(0x00) Error logging NOT supported.
No General Purpose Logging support.

SMART Attributes Data Structure revision number: 5
Vendor Specific SMART Attributes with Thresholds:
Warning: device does not support Error Logging
IOCTL_IDE_PASS_THROUGH does not work on your version of Windows
Error SMART Error Log Read failed: Function not implemented
Smartctl: SMART Error Log Read Failed
Warning: device does not support Self Test Logging
Error SMART Error Self-Test Log Read failed: Function not implemented
Smartctl: SMART Self Test Log Read Failed
Device does not support Selective Self Tests/Logging

Figure 1. Sample Output from smartctl.exe
2.2 Sensing Sound

Three base technologies are needed to support the detection of hardware faults using sound: sensing, or the conversion of sound into electrical signals; digitization, or the conversion of signals into digital data, for use in digital computers; and signal processing, or the analysis of a signal to detect and distinguish its features.

2.2.1 Accelerometers and Microphones

Currently, two types of hardware are commonly used to sense sound: accelerometers and microphones. Accelerometers, which measure mechanical forces in an object due to acceleration, can be used to detect vibration.\(^\text{14}\) Some accelerometers can measure frequencies as high as 30 kHz.\(^\text{15}\) An example of an accelerometer application is the knock sensor used in most modern automobiles to allow the engine control computer to respond to detonation in the combustion chamber. The accelerometer is attached to the engine block and responds to vibrations that indicate the presence of detonation.\(^\text{16}\) While attaching the sensor directly to the engine block reduces the detection of noise from other sources, it does not eliminate all noise and the processing system must deal with this extraneous noise.

Another application for accelerometers has been in detecting vibration from a helicopter transmission. Accelerometers attached near critical points such as bearings and gears produce signals of the vibrations that are analyzed for possible failures.\(^\text{17}\)

An accelerometer attached to a hard drive may detect potential problems with the drive, but miss sounds from other parts of the computer, such as fans, power supplies, or individual electrical components that may emit sounds. Equipping every component with an accelerometer

and other concomitant hardware would add to the cost of the system and the required processing of multiple signal streams would add to computational complexity.

A solution to this problem of multiple sound sources might be to place a microphone inside the computer case to allow it to listen to all the sounds coming from the computer. This, however, presents a different set of problems. A first is ensuring that the microphone is oriented to capture sounds from devices likely to produce sounds of interest. This problem arises with microphones that have a directional component to their response. For example, microphones with unidirectional or cardioid response patterns\textsuperscript{18} are most attuned to sounds at the front end with the response diminishing toward the sides and rear of the microphone.

If a microphone captures sound from every noise-generating component in a computer cabinet, the processing software must distinguish sounds from different sources—a more complicated problem than the case of accelerometers attached to individual components. Further difficulties arise when environmental noises generated from outside the computer cabinet are captured. The processing software must discount these environmental sounds.

2.2.2 Sampling Sounds

In order to analyze the sounds detected by a microphone, the analog voltage produced by the microphone must be converted into a digital value representing the analog signal, a process known as sampling. A device that samples voltage is called an analog to digital converter (ADC). The ADC first captures the input voltage using a process known as “sample and hold.” This fixes the input voltage so that the ADC circuitry can stabilize long enough to do the conversion. The sample represents a snapshot of the signal at a specific moment in time since subsequent changes in the input voltage are ignored. The ADC then pigeonholes the voltage sample into one of a series of possible stepped values that the ADC can accept. This process is known as “quantization.” In this operation, if the sample lies between two permissible digital values, the ADC will choose one of the values and assign it to that value. Then, in the final step,

\textsuperscript{18} Sony Model C-22 FET Condenser Microphone Instruction Manual, Sony Corporation

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the ADC finds a digital number that can represent the quantized voltage.\textsuperscript{19} Quantization loses some information in the conversion, but the loss is unavoidable since a digital number, unlike an analog quantity, can only represent a finite number of possible values. Sampling is repeated at specified intervals at a hardware-determined rate, known as the sampling rate, and typically measured in units of samples/second or Hertz.

Some characteristics of an analog signal may be lost if they occur between samples. This is a result of components of the signal changing faster than the ADC is sampling, i.e., one of the input frequencies is higher than the sampling rate can detect. Nyquist’s Law dictates that the highest frequency that can be reliably represented with a series of samples is half that of the sampling rate. Higher frequencies in analog signals are either lost or appear as lower frequency components.\textsuperscript{20} This phenomenon, known as aliasing, sounds like a metallic clipping in an audio signal.\textsuperscript{21} Signal samples include noise in addition to the desired signals. Noise frequencies should also be considered in choosing a sampling rate, since these may also appear as lower frequency components that could be confused with legitimate signals.\textsuperscript{22}

Increasing the sampling rate solves the problem of undersampled signals, but at a cost of more expensive DAQ equipment to handle a higher sampling rate, a higher demand on the processor, and increased memory requirements. Depending upon the frequencies needed for analysis, sampling at a rate higher than the Nyquist frequency may be a waste of system resources and may not yield additional data meaningful to the analysis.

2.2.3 Sampling Programming Practice

Many contemporary computers come with sound sampling hardware built in. If not, a commercially available adapter can be added. These sound devices, commonly referred to as sound cards, may do a variety of sound-related tasks including digital sampling, playback, MIDI

\begin{flushleft}
\textsuperscript{20} Grover & Deller, 96 \\
\textsuperscript{21} Daubechies, Ingrid, Ten Lectures on Wavelets, Society for Industrial and Applied Mathematics, (1992), 20 \\
\textsuperscript{22} Grover & Deller, 96
\end{flushleft}
synthesis, and sound mixing. These functions are typically available using monaural or multi-channel signals. Sound cards can sample audio signals without extra equipment other than a microphone.

There are constraints on the use of sound card hardware for monitoring sound. One constraint is that sound hardware installed in computers operates in the range of human hearing. This may pose a problem if frequencies outside this range are needed to diagnose a potential failure. Another potential problem is that differences in sound cards may affect the diagnostic tool’s performance when a card is used to monitor hardware.

Software is required to access data captured by the sound card hardware. In multitasking operating systems such as Microsoft Windows or Linux, the operating system manages the sound card, handling hardware interrupts and controlling access to the hardware by application processes. The UNIX/Linux environment supports two standard application programming interfaces (API) for managing sound cards. The now-deprecated Open Sound System (OSS) manages sound hardware using the devices-as-files paradigm typical of the UNIX environment. Under OSS, reading from /dev/audio or /dev/dsp returns sampled data from a properly initialized sampling device.23 The newer, Advanced Linux Sound Architecture (ALSA) supports access to sound cards through calls to the ALSA driver.24 ALSA is the standard sound system in Linux kernel versions 2.6 and higher.

ALSA supports the use of callbacks for data access. A callback is a data-processing function that an application registers with the ALSA driver. When a callback is in use, the ALSA driver accumulates a specified number of sample frames in a buffer that is known to the callback,

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23 The Linux Sound HOWTO, http://www.faqs.org/docs/Linux-HOWTO/Sound-HOWTO.html#AEN611 (accessed 1/11/05)
24 Matthias Nagorni, ALSA 0.9.0 Howto v.0.0.6, http://www.suse.de/~mana/alsa090_howto.html (accessed 2/25/2005)
then redirects processing to the callback. Callbacks reduce processing overhead by eliminating the need for the application to actively monitor (poll) a sound card for incoming data.\textsuperscript{25}

2.3 Analyzing Sounds

2.3.1 Fourier Transform

One of the oldest and most widely used methods of analyzing signals and categorizing them is the Fourier function, developed by Joseph Fourier in 1807. Fourier showed that every continuous periodic function, or signal, can be represented by a sum of sines and cosines as in Equation 1.\textsuperscript{26}

\[
a_0 + \sum_{k=1}^{\infty} (a_k \cos kx + b_k \sin kx)
\]

Substituting different values for the coefficients \(a_k\) and \(b_k\) in Equation 1 yields different functions. Conversely, for any differentiable function \(f\), the \(a_k\) and \(b_k\) that yield \(f\) is determined using an algorithm known as the Fourier Transform. The result is a transformation of a time-domain signal to a frequency-domain representation of that signal.\textsuperscript{27}

Computing the Fourier Transform for a given signal is a non-trivial operation. For a sample size \(n\), the Fourier Transform algorithm involves multiplying an \(n \times n\) matrix resulting in \(n^2\) operations. The algorithmic complexity of the Fourier Transform makes it impractical for real-time signal processing.\textsuperscript{28}

\begin{thebibliography}{9}
\bibitem{25} Paul Davis, \textit{A Tutorial on using the ALSA Audio API} http://equalarea.com/paul/alsa-audio.html#interruptex (2002)
\bibitem{26} Amara Graps, \textit{An Introduction to Wavelets}, Institute of Electrical and Electronics Engineers, Inc., (1995), 2
\bibitem{27} Graps, 5
\bibitem{28} Graps, 5
\end{thebibliography}
2.3.2 Fast Fourier Transform

The Fast Fourier Transform (FFT), a faster algorithm for finding Fourier coefficients developed by Cooley and Tukey,\textsuperscript{29, 30} can be applied to samples obtained at evenly spaced intervals in time. The FFT algorithm uses the assumption of a fixed sampling rate to compute a Fourier matrix product as a product of a few sparse matrices. The resulting algorithm has complexity of $O(n \log_2(n))$, which is fast enough for practical signal processing.

An example use of the FFT is presented in Figure 2 and Figure 3. Figure 2 shows a plot of an input signal created in the mathematical analysis tool Maple using frequencies of 440, 880, 1320, and 1760 Hz., with amplitude values of 1, 0.75, 0.4, and 0.1 respectively.\textsuperscript{31}

![Input Signal](image)

Figure 2. Input Signal

The signal shown in Figure 2 was analyzed using Maple’s FFT function, producing the plot in Figure 3. Spikes occur at positions along the x-axis corresponding to the frequencies and magnitudes used to create the original signal.

\textsuperscript{29} Grover & Deller 356
\textsuperscript{31} Paul Goossens, \textit{Spectral Analysis using Maple 6}, Maple worksheet, Waterloo Maple Inc.
A significant weakness of the Fourier transform is its inability to accurately represent transient components of a signal.\(^\text{32}\) This lack of accuracy stems from the use of the periodic sine and cosine functions to represent aperiodic signals. An extreme illustration of a Fourier representation of a transient signal is presented in Figure 5, which shows the frequency domain of the non-periodic pulse from Figure 4. Unlike the discrete frequencies generated by applying the FFT to the example in Figure 2, applying an FFT to a pulse yields a broad, continuous range of frequencies. The reduced scale in the y-axis indicates that the energy of the pulse is being distributed over the range with much lower amplitude. This use of all frequencies makes FFT representations of pulses impractical. Because many of the sounds that a computer makes are transient, a naive application of Fourier transform may yield inaccurate characterizations of possible syndromes.

\(^{32}\) Graps, 5
2.3.3 Localizing in Time

One strategy for using FFTs to handle transients is to process selected sections of signal at a time. This strategy, the Windowed Fourier Transform (WFT), partitions a signal into windows and then applies the FFT to each window. Transients in the signal can be localized to somewhere within one of these windows. Errors resulting from sharp transitions at the boundaries of a window can be minimized by weighting the window’s center portion.\textsuperscript{33}

\textsuperscript{33} Graps, 5
The WFT, unfortunately, cannot localize transients within windows—a problem when the transients themselves represent the “interesting” part of the signal.\textsuperscript{34} Another issue with the WFT is that the resolution of the window is constant for all frequencies. Small windows are best suited to localizing short events more accurately. On the other hand, larger windows handle low frequencies well, but cannot localize short, high frequency events as well as the small windows. A single window size for all signals may not handle those signals as well as a scheme that uses varying window sizes for different signal frequencies.\textsuperscript{35}

2.3.4 Wavelets

Another technique for representing signals in both the time and frequency domains is wavelet theory. Wavelets are mathematical functions exhibiting specific characteristics that make them particularly useful for analyzing transient signals. One desirable feature of a wavelet function is that it should exhibit compact support—that is, return non-zero values only over a finite domain. The function can be scaled with appropriately chosen coefficients to approximate an input signal and since the wavelet only exists within a specified range, it can be used to approximate transient signals. Wavelet coefficients can be chosen to approximate either large or small windows, thereby allowing the wavelet to localize short-interval, high-frequency events or to handle low frequencies.\textsuperscript{36}

There are two general types of wavelet transforms, continuous and discrete. The continuous wavelet transform is generally represented by Equation 2.
\[
(T^\text{wav} f)(a,b) = |a|^{\frac{1}{2}} \int f(t) \psi \left( \frac{t-b}{a} \right) dt
\]  

(2)

In the continuous wavelet shown in Equation 2, the coefficients \(a\) and \(b\) represent the scaling and translation of a wavelet function, \(\psi\), to allow it to approximate an input signal. The other version of the wavelet transform, the discrete wavelet transform, is shown by Equation 3.

\[
T^\text{wav}_{m,n} (f) = a_0^{-m/2} \int f(t) \psi \left( a_0^{-m} t - nb_0 \right) dt
\]  

(3)

In the discrete wavelet transform shown in Equation 3, the values \(a_0\) and \(b_0\) assume discrete values. There are two types of discrete wavelet transforms: transforms that use frames, non-independent sets of vectors, and transforms in which \(\psi\) and \(f(t)\) are carefully chosen to be orthogonal. In the case of orthogonal transforms, the wavelet, \(\psi\), is scaled by the scaling function \(f(t)\). Sometimes \(f(t)\) is designated by \(\phi\) or \(\varphi\). The family of wavelet functions \(\psi^{a,b}\) is sometimes referred to as the mother wavelet or analyzing wavelet.

2.3.5 Haar Wavelet

One of the first wavelets was the Haar wavelet, described in 1910 by Hungarian mathematician Alfréd Haar. Unlike the Fourier transform, which is a linear combination of sines and cosines, the Haar wavelet consists of a simple step function varying between -1 to 1 over a range of 0 to 1, with the function having a value of zero otherwise. The Haar function is shown in Equation 4 and depicted in Figure 6.

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37 Daubechies, 3
38 ibid
39 Graps, 5

28
Figure 6. Haar Wavelet

Associated with the Haar wavelet function $\psi$ is a scaling function, $\phi$, shown in Equation 5 and depicted in Figure 7.\textsuperscript{40}

$$\phi(x) = \begin{cases} 1 & 0 \leq x < 1, \\ 0 & otherwise \end{cases}$$

\textsuperscript{40} Daubechies, 137
Outside of $[0,1]$, the Haar wavelet is zero, exhibiting compact support. Orthogonality holds when the inner product of the wavelet function and the scaling function is 0, as determined by evaluating Equation 6.\(^{41}\)

\[
\int_{-\infty}^{\infty} \phi(x)\psi(x)\,dx = 0
\] (6)

The Haar wavelet has deficiencies that make it less useful for time localization.\(^{42}\) The wavelet, however, is simple and easy to program. It has been put to many uses, including image compression and the teaching of wavelet theory.

2.3.6 Daubechies Wavelet

Other wavelets have been proposed in addition to the Haar wavelet, including the Mexican Hat,\(^{43}\) Sinc, Morlet,\(^{44}\) and Gabor.\(^{45}\) Yet another family of wavelets, proposed by Ingrid Daubechies, exhibits compact support and orthogonality and is useful as a filter and for time localization of signals. The Daubechies wavelet is defined by a recursive function, shown in Equation 7.

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\(^{42}\) Daubechies, 10
\(^{43}\) Daubechies, 75
\(^{44}\) Daubechies, 76
\(^{45}\) Daubechies, 84
\[
\phi(x) = \sum_{k=-\infty}^{\infty} p_k \phi(2x - k)
\]  \hspace{1cm} (7)

The various Daubechies wavelets differ in the number of non-zero terms. The Daubechies wavelet of interest here, sometimes called the D4 because it has 4 non-zero terms or vanishing moments, can be determined by recursively evaluating Equation 8 over an arbitrary number of iterations, \( n \).

\[
\phi_n(x) = P_0 \phi_{n-1}(2x) + P_1 \phi_{n-1}(2x - 1) + P_2 \phi_{n-1}(2x - 2) + P_3 \phi_{n-1}(2x - 3)
\]  \hspace{1cm} (8)

where the initial condition is defined by \( \phi_0(x) = \begin{cases} 1 & \text{if } 0 \leq x < 1 \\ 0 & \text{otherwise} \end{cases} \).

The refining coefficients \( P_0 \) through \( P_3 \) for Equation 8 are\(^{46} \)

\[
P_0 = \frac{1 + \sqrt{3}}{4} \hspace{1cm} P_1 = \frac{3 + \sqrt{3}}{4} \hspace{1cm} P_2 = \frac{3 - \sqrt{3}}{4} \hspace{1cm} P_3 = \frac{1 - \sqrt{3}}{4}.
\]

When Equation 8 is evaluated over an arbitrary number of iterations, the scaling function \( \phi(x) \) emerges, as shown in Figure 8. The mother wavelet function \( \psi(x) \) can be computed from the scaling wavelet function \( \phi(x) \) using Equation 9, with the resulting wavelet shown in Figure 9.

\[
\psi_k(x) = \sum_{k=1}^{N} (-1)^k \phi_{n-k}(2x - k)
\]  \hspace{1cm} (9)

\[^{46}\text{Daubechies, 235}\]
The D4 wavelet function is neither smooth nor symmetrical and has a fractal self-similarity characteristic. This wavelet $\psi$ can be used to analyze signals by finding coefficients $C_{n,k}$ satisfying Equation 10.

$$ f(t) = \sum_{n=-\infty}^{\infty} \sum_{k=-\infty}^{\infty} C_{n,k} \psi(2^n t - k) $$

Since the wavelet is compactly supported, it is not necessary to sum from $-\infty$ to $\infty$: the wavelet is zero beyond the interval [-1, 3]. Once determined, the scaling coefficients $C_{n,k}$ may be used to reconstruct the original signal by scaling and translating $\phi$ to obtain a good approximation to the signal being analyzed. The coefficients may be manipulated in various ways before reconstructing the signal to provide filtering of the input signal or signal compression. In this research, the coefficients were only used to characterize the input signal; the reconstruction of the original signal was unnecessary.

As an example of signal analysis with the D4 wavelet, consider an analysis of the function $\sin(\pi t)$ where $0 < t < 1$ and zero everywhere else, or $f(t) = \begin{cases} \sin(\pi t) & 0 < t < 1 \\ 0 & \text{otherwise} \end{cases}$.

Using the Daubechies D4 transform, the curve can be approximated by finding values for the

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48 Edward Aboufadel, Daubechies Wavelet, Maple worksheet, 2001, Waterloo Maple, Inc.
scaling and translation coefficients $a$ and $b$. The resulting approximation taken from a Maple Worksheet using just a few coefficients is presented in Figure 10.$^{49}$

In the example shown, the Maple worksheet computes the coefficients using numerical integration. This technique demonstrates the wavelet principle but is too slow and complex to use as a filtering system.

2.3.7 Multiresolution Analysis and the Fast Wavelet Transform (FWT)

A faster way of finding coefficients of a discrete wavelet is Multiresolution Analysis. Multiresolution Analysis finds wavelet coefficients by successively downsampling the signal to lower resolutions. At each stage, the signal is represented by half as many samples as in the previous stage, filtering out high frequencies that cannot be represented at the lower resolution.$^{50}$

$^{49}$ Aboufadel
$^{50}$ Daubechies, 129
This principle can be illustrated by representing a simple sine signal with Haar wavelets. Figure 11 represents a sine function and Figure 12 represents the sine function downsampled using the Haar wavelet, with seven iterations of multiresolution analysis. These figures may suggest that all that has been done is to sample the signal at a lower resolution, and that is partially true. However, re-sampling a signal at a lower resolution or sampling density introduces noise, as determined by Nyquist’s Law, which dictates that a signal must be sampled at twice the highest frequency represented in the signal. Sampling at a lower rate, or downsampling by throwing away every other datum, would introduce aliasing noise into the signal. If high frequencies are removed from a signal by downsampling, the high frequencies that cannot be represented at the lower resolution must first be filtered from the signal to avoid introducing aliasing noise.

At each step in the Multiresolution analysis, the signal is filtered, removing high frequencies that would produce aliasing in the downsampled signal. However, these high frequencies contain information about time-localized events, so it is desirable to preserve this information separately. The input signal can be separated into a high-frequency component representing detail and a low-frequency component to be processed further in the next iteration.
This technique is also known as a Quadrature Mirror Filter or QMF.\textsuperscript{51} Figure 13 depicts the Multiresolution Analysis filter. $V_0$ represents a vector of lower level elements, $V_1$ represents the next level with the detail information removed, and $W_1$ represents the detail information that was removed from $V_1$. The down arrows represent downsampling. The elements $V_1$ and $W_1$ now contain half as many elements as $V_0$ due to the downsampling.

![Figure 13. Multiresolution Analysis](image)

The QMF is applied iteratively to the signal, separating high frequencies into a signal representing high-frequency time events and a signal containing lower frequencies. The process is repeated on the $V_x$ vector as often as needed or until one datum is left. The one remaining element should approximate the average of the sample.\textsuperscript{52} This Multiresolution analysis process is sometimes referred to as the pyramid algorithm or ladder algorithm. illustrates the process of iteratively applying the method depicted in Figure 13. At each level the signal is filtered with high-pass data saved as the detail vector $W_x$; low pass data is then processed in the next iteration, as illustrated in Figure 13.

\textsuperscript{51} Daubechies, 161
\textsuperscript{52} Wim Sweldens and Peter Schröder, \textit{Building Your Own Wavelets at Home}, In Wavelets in Computer Graphics,. ACM SIGGRAPH Course notes (1996), 20
If $V_0$ contains $N$ elements, then both $V_1$ and $W_1$ will consist of $\frac{N}{2}$ elements. Since every level has half as many elements as the level above it, the total number of levels is $\log_2(N)$.

A restriction on the size of $N$ is that it must be a power of 2 or $N = 2^{\text{levels}}$.

The multiresolution algorithm processes half of the input values at each level. The first level processes $\frac{N}{2}$ elements, the next level processes half of the remaining elements or $\frac{N}{4}$, and so on. The total number of operations to complete the transform is $\sum_{k=1}^{\text{levels}} \frac{N}{2^k}$ or an algorithmic complexity of order $O(N)$, a significant improvement over the FFT's complexity of $O(n \log_2(n))$.

A characteristic of Multiresolution Analysis is that the frequencies represented at each level do not evenly subdivide the frequency space. For example, if a signal is sampled at 44,100 Hz, the sampling rate for music CDs, then the sampling rate of the next level will be half the original rate or 22,050 Hz. Each subsequent level will have a sampling rate half the rate of the level before. The frequency subdivision is not evenly spaced as it is in FFT. This is a tradeoff
made in signal analysis. At the lower levels, time localization is good at the expense of frequency resolution. At the upper levels, frequency resolution is good, but there is little time localization. This relationship has been compared to Heisenberg’s Uncertainty Principle, suggesting that we can localize time or resolve frequencies but not both.\textsuperscript{53} Wavelet-based Multiresolution Analysis offers a useful compromise to the problem of analyzing in both time and frequency domains.

### 2.3.8 Lifting

Historically, wavelets have been constructed based upon Fourier series. The wavelets built in this way are termed first generation wavelets. A new way of constructing wavelets has been proposed that allows a spatial construction without Fourier analysis. These wavelets are termed second generation and are built using a method referred to as lifting. The basic idea is to take advantage of the correlated structure of the signal.\textsuperscript{54}

A signal consists of a stream of ordered data that is divisible into even and odd numbered elements. The expected high degree of correlation between the odd and even sets of elements allows either set to predict the values in the other, using prediction functions of the form:

\[ d = x_{\text{even}} - P(x_{\text{odd}}) \]

This is reversible, since from \( d \) and an odd value the even value can be recovered:

\[ x_{\text{even}} = P(x_{\text{odd}}) + d \]

If subsampling is accomplished by selecting the even elements, aliasing will occur. To eliminate this problem, the even values are replaced with a smoothed value using an update operator:

\[ s = x_{\text{even}} + U(d) \]

which can be easily inverted by:

\[ x_{\text{even}} = s - U(d) \]

\textsuperscript{53} Björn Jawerth and Wim Sweldens, \textit{An overview of Wavelet Based Multiresolution Analysis}, SIAM Rev. 36 (1994), no. 3, 377412

\textsuperscript{54} Ingrid Daubechies and Wim Sweldens, Factoring Wavelet Transforms into Lifting Steps, 1996, 2
and then $x_{\text{odd}}$ is recovered from $d$ and $x_{\text{even}}$, as above.\(^{55}\)

The Haar wavelet can be computed with lifting using just two statements of C code where $a$ represents even and $b$ represents odd elements:

\[
b - = a; \quad a += b/2; \quad ^{56}
\]

The inverse is easily accomplished by reversing the order and changing the operators:

\[
a -= b/2; \quad b += a; \quad ^{57}
\]

These operations overwrite the locations that contain the original data with the transformed values, thereby eliminating the need for temporary storage. The simple update operator in this example, an arithmetic average of two values, reflects the simplicity of the Haar transform. More complex wavelets, such as the D4, require more complicated factoring in order to use lifting.\(^{58}\)

### 2.3.9 Wavelet Packets

Dr. Ronald A. Coifman extended wavelet theory by treating the wavelet transform as a special case of the Wavelet Packet Transform.\(^{59}\) The fast wavelet transform described in §2.3.7 and presented in Figure 14 produces a partial tree, obtained by transforming the low-pass filtered component at each level and leaving the high-pass component as output. Coifman’s algorithm transforms the high-pass component as well, thereby generating a full binary tree, illustrated in Figure 15.\(^{60}\) Coifman claims that this transformation provides a more general solution of wavelets, providing more information about the signal. Since there are more calculations required to construct this tree, this transform loses the $O(n)$ algorithmic complexity offered by the Fast Wavelet Transform or Multiresolution Analysis.

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\(^{55}\) Ingrid Daubechies and Wim Sweldens, 1-3
\(^{56}\) Wim Sweldens and Peter Schröder
\(^{57}\) ibid
\(^{58}\) Ingrid Daubechies and Wim Sweldens, 15-16
\(^{59}\) Graps, 10
2.4 Classifying Signals

Procedures have also been developed for classifying signals by comparing their wavelet coefficients to coefficients from known signals.\textsuperscript{61} These procedures include the ANOVA F-test,\textsuperscript{62} neural networks,\textsuperscript{63} and dot product.\textsuperscript{64} Possible solutions might also be found in standard statistical techniques such as logistic regression or chi-square.

2.5 Wavelet Software Solutions

The wavelet software applications listed below are available on the Internet and from other sources. Unfortunately, none of these codes could be used in this project. Some, like the math software plug-ins, could not be used to develop a stand-alone application. Other wavelet software packages were distributed in binary, included source code that would not have adapted well to this project, or were unavailable due to licensing issues.

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\textsuperscript{61} Daubechies, 55
• Rice University provides the *Rice Wavelet Toolbox* (RWT) for Matlab ([http://www-dsp.rice.edu/software/rwt.shtml](http://www-dsp.rice.edu/software/rwt.shtml)). RWT, which interfaces with Matlab, provides 1D and 2D wavelet and filter bank design, analysis, and processing. Matlab ([http://www.mathworks.com](http://www.mathworks.com)) is required in order to use RWT. Neither Matlab nor RWT was evaluated for this project.

• *Wavelet Explorer*, from Wolfram Research, Inc. ([http://www.wolfram.com](http://www.wolfram.com)), runs only with Mathematica, another product of Wolfram Research, Inc. Neither Mathematica nor Wavelet Explorer was evaluated for this project.

• The Maple applications center includes a Maple-based implementation of the D4 wavelet by Edward Aboudafel and an accompanying worksheet ([http://www.maplesoft.com/applications/app_center_view.aspx?AID=965](http://www.maplesoft.com/applications/app_center_view.aspx?AID=965)). Both proved useful for this project, but not for the prototype proper. No other wavelet toolboxes were found for Maple.

• Bear Products International ([http://www.bearcave.com](http://www.bearcave.com)) offers software products by Ian and Linda Kaplan. Their web site includes wavelet articles and source code that can be downloaded free of charge. Unfortunately, their implementations of wavelets appeared to be incomplete. Though the articles were useful, other more direct literature sources were used for this paper.

• *Liftpack* is a wavelet software package for processing wavelets using the lifting scheme written by Gabriel Fernández, Senthil Periaswamy, and Wim Sweldens. Since its primary focus is on image processing, it was deemed unsuitable for this project. Source code can be downloaded from [http://www.cs.dartmouth.edu/~sp/liftpack](http://www.cs.dartmouth.edu/~sp/liftpack).

• *X-Window Wavelet Packet Laboratory (XWPL)* is available for download from [http://math.yale.edu/pub/wavelets/software/xwpl/html/xwpl.html](http://math.yale.edu/pub/wavelets/software/xwpl/html/xwpl.html). The web site notes that XWPL source code will not be released.

• *Wavelets with Integer Lifting* ([http://www.cs.kuleuven.ac.be/~wavelets](http://www.cs.kuleuven.ac.be/~wavelets)) is published under the GNU General Public License. The package, which appears to focus on
image analysis, did not lend itself easily to adaptation to this project. The license included with the C++ source code restricts its distribution for uses other than research.


- **Daubwave** is a C program for wavelet analysis using the Daubechies system of wavelets written by Steven Gollmer. It can be found at several sites online including PC Magazine’s shareware library ([http://shareware.pcmag.com/product.php%5Bid%5D45149%5Bcid%5D93%5BSiteID%5Dpcmag](http://shareware.pcmag.com/product.php%5Bid%5D45149%5Bcid%5D93%5BSiteID%5Dpcmag)). It supports forward and inverse wavelet transforms on signals using several versions of the Daubechies wavelets, including the D4, which can be user selected at runtime. The program, which was designed as a file-based console interface application, would have required modification for this project. The license requires the author’s permission before distributing modified code, which would have restricted its use in this project.

- **ImageLib**, a library of C++ functions for image processing, includes some wavelet software. It can be downloaded from [http://dsp7.ee.uct.ac.za/~brendt/srcdist](http://dsp7.ee.uct.ac.za/~brendt/srcdist) and is free under the GNU GPL. It did not appear to lend itself readily to use in this project because of its emphasis on image processing.

- Amara Graps’ website ([http://www.amara.com/current/wavesoft.html](http://www.amara.com/current/wavesoft.html)) lists additional wavelet software. The site includes wavelet articles and links to other wavelet-related web sites ([http://www.amara.com/current/wavelet.html](http://www.amara.com/current/wavelet.html)), including some of those listed here.

- Another source of information on wavelets, The Wavelet Digest website ([http://www.wavelet.org](http://www.wavelet.org)) offers an archive of past wavelet articles, software, and
messages. A search of the archive and of past messages only revealed the software site of the BearCave (listed above).

2.6 Summary

While the Fast Fourier Transform is a powerful tool for signal processing, there are better tools available for processing signals containing transients. Wavelets preserve time domain information in the signal, but the Haar wavelet does not filter well. The better choice of wavelet is the Daubechies D4. The D4 wavelet offers a compromise between good filtering performance, computational efficiency, and ease of implementation. The Fast Wavelet Transform is a more efficient method of finding coefficients for the discrete wavelet and executes in O(n) time, an important advantage. Finally, lifting should offer a useful advantage by eliminating the need for allocating new storage space for the transform results.
CHAPTER 3

METHODOLOGY

This chapter discusses the development of the prototype software, resources available to the project, limitations and constraints on the results, consistent method of evaluating the success of the prototype, and licensing issues.

3.1 Software Specification

The objective of the project was to develop a software application, using standard PC hardware, that monitors a PC for sounds that indicate abnormal hardware operation, and reports these problems to the user, while imposing a minimal computing load on the host system.

Since the ultimate target was to be a Beowulf Cluster, the application needed to run on the Linux operating system. In order to maintain currency, the system used the ALSA API, the most recent Linux sound API. For efficiency, the system was designed to be interrupt-driven using the ALSA callback feature. In order to reduce costs, the system used a standard PC sound card, requiring only the addition of a microphone to monitor the hardware.

3.2 Programming Environment

A Hewlett-Packard Pavilion 6535 with Gentoo 2005.0, incorporating the Linux kernel version 2.6.11, was used as the development platform and for software testing. ALSA version 1.0.8 was installed with the Gentoo distribution, and the built-in sound card tested for proper operation. This machine was used to prototype the application and to test the prototype. The GNU GCC compiler, version 3.3.5, was used for development under Linux. Some development and testing work was also done in Windows using Microsoft Visual C++ version 6.0.

3.3 Prototype Development

There were four major stages to this study: the development of a filter module, development of the sound data acquisition scheme, development of a classifier, and verification of the application. After a search of the DSP literature suggested the use of the Daubechies D4 wavelet for filtering sound, work was directed toward finding or developing software to implement D4. From the background investigation on Linux sound APIs, ALSA was determined
to be the preferred system to acquire sound data from the sound card. The search for a wavelet classifier revealed few options, and the chi-square method was chosen.

To verify the prototype, failures were simulated by playing recorded sounds from collected hard drives to the system. The results were then compared to the S.M.A.R.T. data obtained from those hard drives.

### 3.4 Filtering

The filtering algorithm chosen needed to be efficient in order to minimize the processing load on the system. The algorithm also needed to detect transients in the signal since many kinds of hardware produce intermittent noises such as clicks or knocks. The Daubechies D4 wavelet (cf. §2.3.6) was chosen as the filtering technique; wavelet packets (cf. § 2.3.9) were rejected as too complex for this project. A search of the Internet (cf. §2.5) failed to find existing D4 software packages that would be suitable for this project. A lifting factoring scheme was tested but when modeled in Excel the algorithm did not function properly, returning values inconsistent with those expected. It was assumed that the Daubechies and Sweldens algorithm was correct and that the Excel model was incorrectly implemented, but the error could not be found and the algorithm was rejected. A simpler implementation of the D4, based on a variation on the Haar wavelet with lifting, was tested and chosen as the filtering algorithm. Since the application does not require reconstruction of the signal, no inverse transform algorithm was implemented.

A C++-based object-oriented design was chosen for the wavelet module so that the wavelet functions would be completely self-contained and the module could be easily incorporated in other applications on other platforms. Design features of the D4 C++ class are discussed in Chapter 4.

### 3.4.1 Collection of Failed Components

Unserviceable hardware devices were collected for study from the East Tennessee State University Computer and Information Science Department’s “Providing Area Schools with

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Technical Assistance” (PASTA) project. PASTA is a capstone project for undergraduates in the ETSU Computer Science Department’s IT concentration that refurbishes old university computers for distribution to local public schools.

Disk drives quickly became the focus of this study, since most of the PASTA-provided components were hard drives. Happily, hard drives are also a major component of computers and one of the most significant sources of sound. The focus on hard drives was not intended to ignore the importance of other hardware components, but it did limit the scope of the results.

3.4.2 Obtaining Hard Drive Data

The sounds generated by the failed hard drives were recorded in a digital WAV sound file for playback during initial testing of the application. These digital recordings were used to help design the system software and to test the prototype.

The hard drives were also classified by condition to form a basis of comparison for the software application. One test of a hard drive’s condition is to determine whether the operating system or BIOS can recognize the drive. This technique can be useful since an electrically unresponsive drive is considered bad. Another test is to determine whether files can be read on the drives. This technique was rejected because there was no convenient way to test for file corruption due to drive failure. Software is available from drive manufacturers to test drives using data-destructive tests such as the “write all zeros” test. According to the documentation, a full test of this type could take several hours or overnight to complete.66 A third method of determining hard drive condition is to obtain S.M.A.R.T. data from the drives using a S.M.A.R.T. software application. Hard drives that do not support S.M.A.R.T. were excluded from the study since it would have been impossible to apply the S.M.A.R.T. test to those drives. In addition, the S.M.A.R.T. technology has become standard in contemporary hard drives.

The software that was used for S.M.A.R.T. testing was the freely available smartmontools application smartctl (cf. §2.1.2). Hard drives were installed in removable hard drive racks, and mounted into rack-equipped PCs, for easier testing. Hard drives that could not be physically mounted in the hard drive rack were excluded from testing.

3.4.3 Testing of Wavelet Transform Class

Initial development of the wavelet class was done using Microsoft Visual C++ 6.0 with further development work in Linux after a Linux machine became available. The wavelet transform class was tested during development using simulated and actual signals recorded from the collection of hard drives. The initial wavelet transform was constructed in an Excel worksheet so that the iterations could be verified. After experiments with the Excel prototype confirmed the algorithm’s validity, the transform was implemented as a C++ class and implemented as a DLL. The DLL was then linked to a Maple worksheet, and used successfully to transform the Daubechies D4 φ function and display the result in a 3D plot. A console-based test application was developed to test the D4 wavelet class and verify correct function of the member functions. This application was successfully compiled and run, using Microsoft’s VC++ compiler and GNU’s GCC compiler.

In addition to testing constructed functions, a Maple worksheet was adapted to import the WAV sound recordings and display animated plots of the transforms. This demonstration illustrated the kinds of data that might be expected from the hard drives in an actual system.

3.5 Sound Sampling

The ALSA API (cf. §2.2.3) was determined to be the preferred Linux API for accessing the sound card. The prototype application was adapted from an interrupt-driven playback shown in the ALSA tutorial. Since this code was licensed under the GNU GPL, the modified code for this research was also released under the GPL. The prototype was developed incrementally by getting it to display captured sound samples, then by adding the wavelet class, and finally by adding the classifier.
3.6 Classifier

Several classification schemes presented in Section 2.4 were considered for this project. Neural-net-based classification was rejected because the work of training a neural net was considered too complex for the prototype. Authors who suggested the ANOVA F-test\(^{67}\) or dot product\(^{68}\) as wavelet classifiers did not report sufficiently satisfactory results to warrant investigating these approaches. A statistical approach called the logistic regression was also rejected as impractical because of the need to recalculate coefficients each time a change was made to the hardware being monitored.

The classification approach that was chosen was chi-square, abbreviated \(\chi^2\). \(\chi^2\) is a powerful, simple, and easy-to-implement algorithm that allows distributions of categorical data to be compared. \(\chi^2\), unfortunately, may also be computationally less efficient than some other techniques. Further development on the classifier is left as future work.

3.7 Limitations

Thirty-two hard drives were obtained for testing. Three of these were excluded: one that did not support S.M.A.R.T. and two that were too large to fit in the test rack. The statistical validity of testing with \(N = 29\) may be an issue.

The hard drives tested were removed from computers being refurbished rather than being randomly sampled from the general population of all hard drives. This means that the statistical results obtained cannot be directly applied to the population of all hard drives and should only be used to determine the validity of SonMon with these particular drives.

The hard drives collected typically were older drives that were not similar to drives currently in production. Applicability of results obtained from these drives to current drives may also be an issue.

\(^{67}\) Raz and Turetsky  
\(^{68}\) Speech Recognition using Daubechies Wavelets
The development machine was equipped with an Intel Pentium II 466 MHz processor. This machine with this processor sufficed for software development but it is unlikely to be similar to the target Beowulf Cluster machines. A more powerful processor may execute the software more quickly, but this assumption was not tested.

3.8 Licensing

The ALSA code modified for this research was adapted from an ALSA tutorial. This tutorial is licensed under the GNU General Public License. Under the terms of the GPL, modified code may be freely released to the public. Pursuant to the provisions of the GPL, the SonMon code presented in the appendix may be distributed freely under the terms of the GPL.

Unlike the ALSA code fragment, the D4 wavelet class was developed independently by the author. Since it is identified as a separate component of the application, it is not covered by the GPL. However, in the interest of further development of this research and the concept of monitoring computers via sound, the D4 wavelet code may also be freely distributed under the terms of the GPL.

The text of the GNU General Public License can be found in the appendix. The GPL text should be included with any redistribution of modified versions of the application described here.

3.9 Summary

The wavelet transform was developed by modeling the algorithm in Excel, then coding the class in C++. The ALSA code fragment used for this research was adapted from the ALSA tutorial. Since the ALSA tutorial is released under the GNU GPL, the modified code is also released under the GPL. The classifier was developed as an implementation of the $\chi^2$ statistical technique.

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CHAPTER 4
THE APPLICATION

This chapter discusses the software design and operating principles of the SonMon prototype (name derived from the Latin root *Sonus*, meaning sound and the word *monitor*). The first section presents an overview of system operation. The following sections discuss SonMon’s major components, with a description of the module design and explanation of module operation.

4.1 Overview of SonMon

The SonMon prototype, implemented in C++, demonstrates the viability of monitoring PC hardware using sound. The system acquires sound from a microphone connected to the PC’s built-in sound card, transforms the data using the Daubechies D4 wavelet transform in a Multiresolution Analysis algorithm, and compares the resulting set of wavelet transform coefficients to a template.

Figure 16 illustrates SonMon’s operation. The system loops continuously, analyzing signals from the microphone and sending messages when the system detects faults. The comparison algorithm constructs a histogram by counting specific coefficient values, compares the counts in each histogram bin with a template using $\chi^2$ to compute an error score, and sums the $\chi^2$ value over the entire histogram. A message is issued when the $\chi^2$ total exceeds a pre-determined threshold. When used as part of a statistical test, $\chi^2$ test results are converted into a probability score based on the $\chi^2$ distribution. In this application, conversion to a probability is unnecessary, and the $\chi^2$ number alone is used to evaluate the comparison.

4.2 D4 Wavelet Class

The D4 wavelet class was initially developed and tested in Windows using Microsoft VC++ 6.0. The D4 class was verified by compiling it into a Windows DLL for use with a Maple worksheet. The DLL code first instantiates a wavelet object, then passes data from the Maple worksheet to the wavelet, invokes the transform function, and returns the data to Maple in a two-dimensional array. The code for the DLL is included in the appendix.
4.2.1 Description of Code

As described in Section 2.5, a search for existing D4 wavelet software yielded no satisfactory codes. Consequently, wavelet software was developed from the Daubechies formulas. The wavelet software was developed as a C++ class in order to make the wavelet module portable to other platforms. The class has been compiled and tested on both Linux and Windows platforms without modification. Since SonMon does not require an inverse wavelet transform, no inverse function was developed in the D4 class. While lifting as described in Chapter 2 is appealing due to its ability to reuse the input array, the algorithm found in the referenced work produced unexpected results when tested in an Excel worksheet simulation. The algorithm that was used was based on direct calculation using the Daubechies D4 formulas. This version worked correctly when tested in an Excel worksheet giving expected results and

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worked again when tested in Maple. This algorithm also conserves memory by reusing its input array to calculate the transform.

The D4 wavelet class encapsulates the wavelet data, applies the D4 wavelet transform, and provides access to the transformed data. Data used by the transform are held in a 1024-element array of double-precision floating-point variables. This array, which initially contains incoming data, is repeatedly overwritten over the course of the transform. Double-precision floating-point variables are required because the D4 wavelet transform multiplies each element with real coefficients. Because the incoming data and the transformed data use the same array, no additional storage is required to process and store the transform except for a fixed set of temporary variables in the transform function. When a new stream of samples arrives, the data input function accepts the data elements and stores them sequentially in the array maintaining time ordering of the data.

The transform function will not execute until all 1024 elements have been loaded. The function returns 0 on successful completion and -1 on failure. Attempts to transform previously transformed data are ignored, since this would produce undefined results; the function returns 0 instead, to indicate that the already transformed data are still available. The transform function will not execute again until the class’s reset function is called and 1024 elements of new data are loaded.

The transform function stores the result of the transform in the same data array used for input data by interleaving the coefficient results within the array. Subsequent iterations of the transform access appropriate interleaved values and leave coefficient values from previous iterations undisturbed. Interleaving data eliminates the need to sort intermediate coefficients into separate V and W vectors following each iteration of the transform algorithm.

Access to the interleaved transform coefficients is supported by two accessor functions: `getLevel` and `getLevelSq`. Each accessor function takes two integers as arguments. The first, which represents the level of the transform, is a value between 0-9, corresponding to the ten (i.e., \( \log_2{1024} \)) levels obtained by applying the D4 transform to 1024 elements of data. The second
parameter is index of the desired element within that level. The only difference between these functions is that getLevelSq squares the data before returning it.

The D4 wavelet class was tested by code inspection and by incorporating it into a Windows DLL function to allow testing using Maple. A Maple-based test suite was used to test the class’s operation on several mathematical functions, including the D4 $\phi$ function, to ensure that the class returned expected results. A test application written for C++ was also used to test for proper operation of the member functions. This test application included hard-coded data for the D4 $\phi$ function. A D4 transform of $\phi$ should return all 0’s except for one non-zero value at the highest level. The test application tested all accessor member functions by displaying the transform values. The member function that returns the number of values in each level of the transform was also tested demonstrating that it returned the correct value for each specified wavelet level.

4.2.2 Spike Signal Demonstration

To demonstrate the wavelet class, the D4 transform was applied to the spike signal depicted in Figure 4, the results of which are presented in Figure 17. Figure 17 was prepared using Maple linked to the Windows DLL containing the D4 wavelet class. The axis “Level” in Figure 17 indicates the iterations of the Multiresolution analysis where level 0 is the first pass and level 10 is the end of the transform. The X-axis corresponds to the number of elements in each iteration. Although there are 1024 elements in the original signal and 512 elements in level 0 of the transform, only 100 elements are depicted in this plot, with the rest omitted for clarity. The line of spikes is sometimes referred to as a “ridge” and the irregular values as “ribs.” Compared to the FFT transform in Figure 5, this analysis retains distinct information about the original signal.

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4.3 Main Function

The main function initializes the system and then enters an idle loop, leaving the callback function to monitor system sounds. Initialization consists of setting operating parameters for the sound card and registering the ALSA callback function async_callback with ALSA. The system is initialized to allow interleaved operation with only one input channel, data format as little-endian 16-bit signed, and 44,100 samples/second operation. The sampling rate parameter can be changed, but any changes to the sampling rate may also necessitate changes to the detection parameters.

4.4 Callback Function

The callback function, async_callback, is invoked by way of a sound-card interrupt when the sound card accumulates 1024 unprocessed frames of data. The sound-card interrupt is handled by the operating system, which transfers control to ALSA. ALSA, in turn, transfers control to async_callback, which copies data from the ALSA buffer to the wavelet object via the ALSA function snd_read and the wavelet object’s add function. When all 1024 frames have been copied, the wavelet object’s transform function is invoked, performing the D4 wavelet transform.
After the D4 transform function has executed, the callback function calls the \textit{classify} function to classify the transformed data. If the $\chi^2$ value returned by \textit{classify} exceeds a predetermined threshold (cf. line 79, sonmon.c, Appendix B), \textit{async\_callback} sends a fault message to standard output. (In a production version of SonMon, error messages would be dispatched independently of interrupt processing; the callback would set a flag and \textit{main} would send fault messages, during periodic awakenings from \textit{sleep}.) Resettable flags help to ensure that only one message is displayed for each good/bad transition rather than to display a message for each fault.

\section*{4.5 Classify Function}

The \textit{classify} function first constructs a histogram by binning its inputs. Binning is the operation of counting the number of elements that lie within specified values as if the numbers were being pigeonholed into bins. After binning the elements, \textit{classify} then compares the counts in this histogram to a template, yielding a $\chi^2$ score. Finally, \textit{classify} uses this score to rate the match between the input signal and the template.

\subsection*{4.5.1 Histogram}

A histogram is a count of objects that fit a set of specified criteria. The histogram generated by \textit{classify} counts the number of wavelet coefficients at each level that can be assigned to a set of pigeonholes or bins. The histogram is constructed from data returned by the wavelet class’s \textit{getLevelSq} function. The coefficients returned by \textit{getLevelSq} are squared, which simplifies binning by eliminating the need to handle negative numbers. Since most of the data appear to cluster near 0, a non-linear bin arrangement was chosen that uses a first bin that contains values in the range $[0,2^6)$, and ten progressively larger bins, where each bin is twice the size of its predecessor. The twelfth and final bin contains all values above $2^{17}$ or 131,072. This non-linear bin arrangement, which allows the bins to fill more evenly, proved adequate for classifying the data obtained for this research. One area of future work might be to test other possible bin arrangements to determine optimum performance.

The current algorithm for binning uses a linear search on bin ranges to map each (squared) coefficient to its proper bin. This algorithm, while inefficient ($O(n^2)$), was sufficient
for testing the prototype. A more efficient algorithm for binning this data should be used in a production version of SonMon.

Histograms for the wavelet tree’s higher levels may not support useful classifications of input data. A histogram generated from the tree’s top level, which contains one value, is plainly ambiguous. A histogram generated from the tree’s second level, which contains two elements, is marginally more useful. Currently, classify uses the wavelet’s lower 6 levels. These levels contain more high frequency detail, which could be filtered by ignoring those levels. Higher levels of the wavelet contain more thoroughly filtered levels that may contain meaningful information. Future work could determine which levels of the wavelet are most useful.

4.5.2 Chi-Square Comparison

When the histogram has been completed for a single level, a $\chi^2$ comparison is used to evaluate the match between the histogram and a template. The $\chi^2$ results for each comparison are summed and the function returns the total $\chi^2$. The prototype currently compares the lowest six levels of the wavelet (cf. §4.5.1). Another possible problem with the $\chi^2$ algorithm is that it potentially violates the $\chi^2$ rule of thumb that dictates that the minimum count in the template be at least five.

The $\chi^2$ threshold value is hard-coded at 1000 and was initially arrived at by trial and error. Optimization of the threshold value is discussed in Chapter 5.

4.6 Templates

The template used in this work was prepared by modifying the code to output histograms to a text file. The modifications to the code can be found in lines 96, 133-140, 155, 371, and 387 of the sonmon.c code shown in the appendix. These lines may be uncommented and the code recompiled to activate these lines. They produce an output file, D4Output.txt, containing the histogram data as well as the $\chi^2$ value calculated on the current template. The prototype’s template histogram was generated from a representative data block from one of the drives that
S.M.A.R.T. had evaluated as good: one that represented the sound of the drive’s steady state operation.

4.7 Summary

The D4 class is a mature code: one that can be used in final versions of the SonMon application. The main function and the callback function should be modified for use in a production application. The prototype is not suitable for production use but has features to allow its use for further development and proof of concept.

The classify function may need to be modified for further use. Changes to the histogram code may be required and a provision to allow multiple histograms should be added.
CHAPTER 5
RESULTS

5.1 Introduction

SonMon was evaluated by comparing its classifications of 29 hard drives with S.M.A.R.T. data obtained from those drives. These test drives were provided through ETSU’s PASTA program (cf. §3.4.1). Sound recordings were obtained of these drives by connecting them to power but not to a computer. S.M.A.R.T. data then were extracted from these drives by connecting each drive to a computer and running an application to extract the S.M.A.R.T. data. Test results were obtained from SonMon by running it with the sound recordings of the drives. The SonMon test results were compared with the S.M.A.R.T. data to determine whether SonMon had obtained similar results. S.M.A.R.T. was used to determine the condition of the drives and SonMon was evaluated for its ability to agree with the S.M.A.R.T. data.

In addition to S.M.A.R.T. data, digital sound recordings were obtained from each hard drive using Microsoft Sound Recorder. Because some hard drives contained possibly sensitive data and because of the weight, bulk, and delicate nature of the hard drive themselves, they could not be physically transported. Therefore, recordings were taken on a notebook computer for portability. These recordings were used to evaluate SonMon rather than monitoring the hard drives directly. The recordings offered the additional advantage of allowing significant events to be isolated and repeated, a task that would be difficult to do when working directly with the drives.

The results from evaluating the error counts using SonMon were compared to the results obtained from S.M.A.R.T. One test used is Cohen’s kappa (κ), which showed fair agreement between SonMon and S.M.A.R.T. The Binomial Distribution was used to compare the counts of good and bad drives between SonMon and S.M.A.R.T. to test the hypothesis that any agreement was by chance. A p-value was calculated from the binomial to test if the probability of agreement was better than chance. If SonMon is performing as expected, this probability should
be a value less than 5%. However, a less strict 10% confidence level may be justified because of expected inconsistencies between SonMon and S.M.A.R.T. results.

5.2 S.M.A.R.T. Test Procedure

To obtain S.M.A.R.T. data from the sample hard drives, each drive was installed in a PC computer equipped with a removable hard drive rack. The computer was configured so that its internal hard drive was installed on IDE 0 and could boot the machine into either Windows or Linux. So that the drives could operate independently of the main hard drive, the removable hard drive rack containing the suspect drives under test was connected to IDE1, and the drives in each case were jumper-configured as master. With the computer configured in this way, an inoperative hard drive would not keep the machine from booting from the primary hard drive.

The drives were tested with the PC running Windows. The tests were repeated on some drives using Linux, to confirm that drives would not respond to the PC’s system BIOS or Windows had suffered electrical failure. Smartctl version 5.33 was used to attempt to obtain S.M.A.R.T. data from the drive. The smartctl command used was:

    smartctl -a -T verypermissive /dev/hdb

The smartctl -a option retrieves all data from the drive. The -T option with the “verypermissive” setting forces smartctl to ignore S.M.A.R.T. errors. This option was chosen after some drives ignored the default “normal” setting and falsely reported that unimplemented S.M.A.R.T. features were enabled. The smartctl manual reports this as a possible outcome. A drive was considered to have properly responded to S.M.A.R.T. if smartctl displayed the S.M.A.R.T. data section as well as the drive’s information section. A sample smartctl report is shown in Figure 1.

Some drives responded to smartctl with an information section but did not return any S.M.A.R.T. data beyond the drive ID. In these cases, the messages returned from smartctl

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indicated that S.M.A.R.T. was available but not enabled. In these cases the smartctl command “smartctl -s on” was used in an attempt to enable S.M.A.R.T., and then the original smartctl request was repeated. If the drive continued to fail to respond to smartctl with S.M.A.R.T. data, the drive was classified as having failed.

5.3 S.M.A.R.T. Test Evaluation

The results of the S.M.A.R.T. tests are presented in Table 1. Eleven drives were classified as good, including ten that were classified as good by S.M.A.R.T., and one drive that had previously failed a S.M.A.R.T. test (according to a S.M.A.R.T. report), but that was now passing its test and functioning normally. Of the remaining nineteen, eighteen were classified as bad, after failing the S.M.A.R.T test; failing to respond with S.M.A.R.T. data, even after an attempted reenabling of S.M.A.R.T.; or failing to respond at all. A remaining drive could not be tested using available hardware.

<table>
<thead>
<tr>
<th>Test Result</th>
<th>Bad</th>
<th>Good</th>
</tr>
</thead>
<tbody>
<tr>
<td>S.M.A.R.T. reports “passed”</td>
<td>10</td>
<td></td>
</tr>
<tr>
<td>S.M.A.R.T. reports “Passed” with a past history of failure</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>S.M.A.R.T. reports “failed”</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>Drive did not respond with S.M.A.R.T. data</td>
<td>5</td>
<td></td>
</tr>
<tr>
<td>Drive did not respond at all</td>
<td>11</td>
<td></td>
</tr>
<tr>
<td>TOTAL, by category</td>
<td>18</td>
<td>11</td>
</tr>
</tbody>
</table>

5.4 SonMon Test Results

To test SonMon, sound recordings of sample hard drives were played into a microphone connected to a Linux computer. SonMon was tested with recorded sound from the sample hard drives. Recordings were obtained using a Dell 600m computer with a Sony Model C-22 FET condenser microphone connected to the computer’s mic input. Recordings were made using the Microsoft Sound Recorder utility with a sampling rate of 48,000 samples/second and 16-bit monaural samples. A high sampling rate was chosen to maximize fidelity. The microphone was positioned approximately ½” from the left side of the drive, as viewed from the connector end. For consistency, settings were held constant during a recording session. Because the recordings
were made in two separate sessions, and since Windows Volume Control has no convenient way of setting the microphone volume control numerically, it was difficult to keep the volume setting consistent.

Two different computers were used to do the recordings. The Dell 600m computer used to record the first session failed and the second recording session was done with a different system board installed, introducing possible inconsistencies between the recordings in the set. Efforts were made to match the volume between recordings, but some variation may exist.

Rather than analyze the digital sound files directly, the recordings were played from a Windows desktop computer into a Linux machine running SonMon, capturing the sounds using a Sony C-22 FET microphone placed in front of one of the speakers. This arrangement was intended to simulate how the system would operate when monitoring a PC computer. The desktop computer used for sound reproduction was a Pentium II 400 MHz machine using a Tyran S1846SLA system board equipped with Panasonic EAB 710P speakers. The microphone was placed in front of one of the speakers and the Windows volume control was manually adjusted to minimize the error counts displayed by SonMon.

Sound recordings were used instead of monitoring the hard drives directly for reasons of convenience. The SonMon computer was located in a different building from the lab that housed the donor disk drives. Permission could not be obtained to remove hard drives that contained sensitive data from the donor lab intact. Even if such permission had been granted, the collection of hard drives would have been bulky and difficult to move. Recordings facilitated the transfer of data between the labs. The recordings also allowed sound events of interest to be isolated and played repeatedly to verify results.

The $\chi^2$ algorithm used in SonMon is sensitive to the sound input volume. If the sound being sampled exactly matches the sound used to create the template being used, the $\chi^2$ will return a value of zero. Any deviation from a perfect match results in greater values. While testing SonMon, the volume control of the playback machine was adjusted to minimize the number of $\chi^2$ error counts. Once the volume was adjusted, the test was re-run to obtain the test data output file.
Since the SonMon prototype is currently configured to monitor only steady state sounds, Microsoft Sound Recorder was used to extract steady-state sounds from the original recordings, removing the sounds of drive startup and shutdown. Hard drive startup usually includes the sound of the drive spinning up and some head access during initialization. Shutdown sounds include the sound of the power supply switch being turned off and the drive spinning down. Typically, these intervals included the first 5-7 seconds at the beginning of each recording and approximately 10 seconds from the end of each recording. The remaining sounds were representative of steady state operation. Some of the drives exhibited abnormal sounds during this operation and detecting these abnormal sounds is an objective of SonMon.

The volume control for the microphone was set at 100% using alsamixer. Any fine adjustments in sound volume were accomplished at the playback computer using its Windows volume control.

SonMon was modified slightly to write $\chi^2$ scores to a text file by uncommenting lines 96, 155, 371, and 387 of the sonmon.c code shown in the appendix. Compiling SonMon with these lines un-commented causes SonMon to write error values to a file named D4Output.txt in the same directory as SonMon. This code outputs error scores to D4Output.txt until SonMon terminates. The D4Output.txt file was subsequently renamed with hard drive ID numbers for identification and the file imported into Microsoft Excel for evaluation. In Excel, $\chi^2$ values exceeding the threshold of 1000 currently coded into SonMon were counted, and the error counts were tabulated.

Because SonMon produces very high error values and large numbers of error counts when no sound is being received, there were a series of large numbers in the output file at the start until the recording was started. A drop in error values in the input file were taken as the start of the test recording and the point where very high values returned to the no-signal level was taken as the end. Only the values corresponding to the recording were included in the analysis.

When the data were imported into Excel, $\chi^2$ values greater than the threshold value used in SonMon were counted. Counts of 0 were considered good and drives with large $\chi^2$ counts
were considered bad. Drives with non-zero counts close to 0 were classified as suspicious and tabulated separately from the other two groups. Using an initial threshold value of 1,000, these drives typically had error counts ranging between 1 and 20.

Because the length of the input recordings varied between drives, the number of total frames produced by SonMon varied. Since this variation in total frames could affect the evaluation of these suspicious drives, the error count was also calculated as a percent of total for those drives. The total number of frames for each drive averaged 901, with a maximum of 1,357 and a minimum of 583. Because SonMon was sampling at 44,100 Hz. and there were 1,024 points in each frame, 901 frames corresponds to approximately 21 seconds of sampling time with the range of recording times of between 31 and 13 seconds.

Eleven drives were classified as good because they had error counts of 0, indicated by there being no $\chi^2$ values greater than 1,000. Seven drives had error counts greater than 20, indicated by the count of $\chi^2$ values exceeding 1,000. The remaining drives had error counts ranging between 1 and 20, qualifying them as suspicious. Since these drives lie in a boundary, they were tabulated separately. The results of the tests are shown in Table 2.

<table>
<thead>
<tr>
<th>Test Evaluation</th>
<th>Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>Good</td>
<td>11</td>
</tr>
<tr>
<td>Suspicious</td>
<td>11</td>
</tr>
<tr>
<td>Bad</td>
<td>7</td>
</tr>
<tr>
<td>Total</td>
<td>29</td>
</tr>
</tbody>
</table>

5.5 Evaluation of Results

This comparison of S.M.A.R.T. tests with SonMon tests is less than ideal, since the tests assess different indicators of disk drive performance. SonMon cannot identify hard drive failures that produce no characteristic sound: for example, a failure resulting from a broken signal wire or a malfunctioning disk cache chip. The drive may spin up and sound normal but will not function because of the electrical failure. On the other hand, SonMon may detect faults that S.M.A.R.T. misses because they have not yet deteriorated to the point of failure. For example, the one good
hard drive that had a failure incident in its S.M.A.R.T. history (cf. §5.3) exhibits abnormal sounds that were detected by SonMon.

The SonMon template is currently optimized to “listen” for abnormal operation when the drive is spinning and not being accessed. Configuring SonMon to monitor other normal operations like normal head movement during normal drive access, spin up, and shut down will be left for future work. SonMon might not be expected to monitor spin up and shut down when the computer is being started up or shut down, since SonMon may be inactive during these phases of computer operation. If the hard drive is being placed into standby in order to save energy, then SonMon might be used to monitor hard drive spin up or shut down during these periods. SonMon has not been tested with these and other hardware device sounds due to the constraints under which this research was conducted, and the low numbers of sample components available for this research.

Because of these issues and the low data sample, the usual 5% level of confidence for statistical comparison was abandoned, in favor of a less rigorous 10% level of confidence.

All drives with non-zero error counts were classified as bad and drives with error counts of 0 were considered good. The drives that were listed in Table 2 as suspicious were classified as bad for purposes of evaluation of SonMon. This is the only unambiguous partition possible. The evaluations for the S.M.A.R.T. tests and the results from SonMon tests are combined in Table 3.

<table>
<thead>
<tr>
<th>S.M.A.R.T.</th>
<th>SonMon</th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>good</td>
<td>bad</td>
</tr>
<tr>
<td>good</td>
<td>6</td>
<td>5</td>
</tr>
<tr>
<td>bad</td>
<td>5</td>
<td>13</td>
</tr>
</tbody>
</table>

5.5.1 Agreement of Observers Test, Kappa

A common statistical test for quantifying agreement between two observers is Cohen’s kappa (κ). Kappa is a value ranging between 0 and 1 that determines the degree with which two
observers agree on a data set. A kappa of 0 corresponds to an agreement that is no better than random chance while a kappa of 1 indicates perfect agreement.

Kappa for Table 3 was 0.2677, indicating moderate agreement between the evaluations by S.M.A.R.T. and SonMon.

5.5.2 Probabilities of Disagreement

Another statistical test that may be applied to the data is to calculate the probability of disagreements between S.M.A.R.T. and SonMon. Table 3 shows that there were 5 cases in which S.M.A.R.T. indicated that the drive was good and SonMon indicated that the drive was bad. The probability of this case, \( P(B_{SonMon}|G_{SMART}) \) is \( \frac{5}{29} \) or 0.17. Similarly, Table 3 shows that there were also 5 cases in which S.M.A.R.T. showed the drive was bad while SonMon indicated the drive was good. The probability of this case, \( P(G_{SonMon}|B_{S.M.A.R.T.}) \), is also \( \frac{5}{29} \).

5.5.3 Binomial Distribution

The binomial distribution is used as the primary means of determining agreement between S.M.A.R.T. and SonMon. For the binomial distribution test, the null hypothesis \( h_0 \) is that the probability of agreement is by chance or \( \pi=0.5 \). The alternative hypothesis \( h_a \) is that the probability is better than chance or \( \pi > 0.5 \). Since out of 29 drives, there were 19 agreements, the p-value is the sum of the binomial probabilities where \( x \geq 19 \), a calculation which yields a p-value of \( p=0.068 \). This p-value is less than the expected confidence value of 0.10 so \( h_0 \), the hypothesis that the agreement is by chance is rejected. SonMon and S.M.A.R.T. do agree within the expected level of confidence.

5.6 Improving Performance

5.6.1 Templates

The template used in the foregoing tests was a histogram generated from one of the good hard drives (cf. §4.5.1, 4.6). While the system can identify sounds similar to ones in the template, other sound conditions will need to be evaluated before classifying a drive as bad or good.
An example of a different event is the sounds produced by read/write head movement during disk access. These sounds will not match the template used in the tests. A separate template is needed to identify normal disk access as well as any other normal sound that may not be identified by the template used in the SonMon prototype. Furthermore, each installation of the software should have a customized template or set of templates that are unique to that installation. Using a generic template that does not consider individual variations between systems will produce mediocre results.

While the foregoing tests yielded satisfactory results, there was room for improvement in accuracy. The drives that were classified as good by S.M.A.R.T. had an approximately even chance of being classified correctly by SonMon (cf. Table 3). In addition, drives that SonMon classified as good had an approximately even chance of having been similarly classified by S.M.A.R.T.

5.6.2 Threshold Value

Adjusting the threshold value used to evaluate the $\chi^2$ score produced by the classify function could also have improved SonMon’s ability to classify disk drives, relative to the set of disk drives used in this experiment.

Reducing the threshold tends to result in more good drives being classified as marginal or bad, and more marginal drives being classified as bad. For a threshold value of 900, for example, 4 drives are classified as good by both S.M.A.R.T. and SonMon, and 17 drives classified as bad by both. Seven drives classified by S.M.A.R.T. as being good were classified as bad by SonMon and one drive that S.M.A.R.T. detected as bad was passed as good by SonMon. SonMon and S.M.A.R.T. agreed on the evaluation of 21 drives and disagreed on 8 drives. These results are summarized in Table 4.
Table 4. Threshold 900 Results

<table>
<thead>
<tr>
<th></th>
<th>SonMon 900</th>
</tr>
</thead>
<tbody>
<tr>
<td>good</td>
<td>4</td>
</tr>
<tr>
<td>bad</td>
<td>7</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th></th>
<th>S.M.A.R.T.</th>
</tr>
</thead>
<tbody>
<tr>
<td>good</td>
<td>1</td>
</tr>
<tr>
<td>bad</td>
<td>17</td>
</tr>
</tbody>
</table>

Computing kappa for Table 4 yields a value of 0.3446. This value indicates improved agreement between S.M.A.R.T. and SonMon. Computing the p-value for the 21 agreements out of 29 drives, the binomial distribution for \( x \geq 21 \) for this case shows that the p-value has improved to 0.012. This p-value suggests rejecting the \( h_0 \) case, since it is less than the 0.10 level of confidence. The adjustment to the threshold value improves the significance score while slightly increasing the number of false alarms. With the threshold at 900, seven good drives were classified as bad by SonMon, as compared to threshold value of 1000 where only five good drives were incorrectly classified as bad. The big benefit came from the change correctly classifying as bad nearly all the drives that S.M.A.R.T. had classified as bad.

Increasing the threshold shifts results in the other direction. More bad drives are passed as marginal or good, and more marginal drives are classified as good. The result of using a threshold of 1100 is shown in Table 5.

Table 5. S.M.A.R.T. vs. SonMon with Threshold 1100 Results

<table>
<thead>
<tr>
<th></th>
<th>SonMon 1100</th>
</tr>
</thead>
<tbody>
<tr>
<td>good</td>
<td>10</td>
</tr>
<tr>
<td>bad</td>
<td>1</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th></th>
<th>S.M.A.R.T.</th>
</tr>
</thead>
<tbody>
<tr>
<td>good</td>
<td>9</td>
</tr>
<tr>
<td>bad</td>
<td>9</td>
</tr>
</tbody>
</table>

Calculating kappa for Table 5 yields a value of 0.358. This kappa is approximately the same as the kappa for the threshold = 900 case, indicating that agreement is similar. The binomial distribution p-value for this case is the same as for the threshold = 1,000 case, i.e., \( p = \)
0.068. In this case, only one good drive was classified as bad while half the bad drives were passed as good.

To find the optimum value for the threshold, threshold values were evaluated for a range of values from 800 to 1100 at intervals of 10. The results, showing correct and incorrect counts and the associated binomial p-value, are displayed in Table 6. Those p-values less than 0.10, the confidence level chosen for this part of the project, are highlighted in the table to show those values that are lower than the confidence value of 0.10.

<table>
<thead>
<tr>
<th>Threshold</th>
<th>Agree</th>
<th>Disagree</th>
<th>p</th>
</tr>
</thead>
<tbody>
<tr>
<td>800</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>810</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>820</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>830</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>840</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>850</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>860</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>870</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>880</td>
<td>19</td>
<td>10</td>
<td><strong>0.068023</strong></td>
</tr>
<tr>
<td>890</td>
<td>21</td>
<td>8</td>
<td><strong>0.01206</strong></td>
</tr>
<tr>
<td>900</td>
<td>21</td>
<td>8</td>
<td><strong>0.01206</strong></td>
</tr>
<tr>
<td>910</td>
<td>21</td>
<td>8</td>
<td><strong>0.01206</strong></td>
</tr>
<tr>
<td>920</td>
<td>19</td>
<td>10</td>
<td><strong>0.068023</strong></td>
</tr>
<tr>
<td>930</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>940</td>
<td>17</td>
<td>12</td>
<td>0.229129</td>
</tr>
<tr>
<td>950</td>
<td>17</td>
<td>12</td>
<td>0.229129</td>
</tr>
<tr>
<td>960</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>970</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>980</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>990</td>
<td>19</td>
<td>10</td>
<td><strong>0.068023</strong></td>
</tr>
<tr>
<td>1000</td>
<td>19</td>
<td>10</td>
<td><strong>0.068023</strong></td>
</tr>
<tr>
<td>1010</td>
<td>19</td>
<td>10</td>
<td><strong>0.068023</strong></td>
</tr>
<tr>
<td>1020</td>
<td>19</td>
<td>10</td>
<td><strong>0.068023</strong></td>
</tr>
<tr>
<td>1030</td>
<td>20</td>
<td>9</td>
<td><strong>0.030714</strong></td>
</tr>
<tr>
<td>1040</td>
<td>17</td>
<td>12</td>
<td>0.229129</td>
</tr>
<tr>
<td>1050</td>
<td>17</td>
<td>12</td>
<td>0.229129</td>
</tr>
<tr>
<td>1060</td>
<td>17</td>
<td>12</td>
<td>0.229129</td>
</tr>
<tr>
<td>1070</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>1080</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>1090</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
<tr>
<td>1100</td>
<td>18</td>
<td>11</td>
<td>0.132465</td>
</tr>
</tbody>
</table>
From Table 6, it can be seen that the choices of threshold values of 900 and 1000 were fortunate. The program’s response to changes in threshold values appears to be very critical. In view of these results, the choice of threshold value should be made carefully.

5.6.3 Histograms

Another opportunity for improving performance is in optimizing the histogram template used to evaluate the wavelet coefficients. The template used in the current prototype was obtained by modifying the software to export a sample of histogram values to a file. The code to export the histogram is included in the code shown in the appendix. Sonmon.c lines 96, 133-140, 155, 371, and 387 should be un-commented and the application re-compiled in order to activate this feature.

The text file obtained from SonMon with the modifications can be imported into a spreadsheet for study. The objective is to find a histogram frame that corresponds to a sound that represents a normal mode of operation. The histogram currently in use represents a quiescent state with no head activity. After sampling the sounds for a few seconds, SonMon should be terminated with the ctrl-C key. The data will be stored in the D4Output.txt file. To build a histogram, the numbers are copied from the chosen histogram block into the array declaration at the head of the application and the application is re-compiled.

A histogram template should yield a $\chi^2$ value of 0 when a sound identical to the sound used to create the template is encountered. No values smaller than 0 can exist but mismatches between the test sound and the template will result in positive values. Normal sounds may not always match the template and additional templates may be needed to match the histogram. SonMon needs a feature added to allow the successive matching of multiple histograms. An example of the need for multiple histograms is the normal sounds of read/write head access. A histogram template that matches the sound of the spinning platter will probably classify the sounds of normal head access as an error. A second histogram, optimized for head access sounds, could identify this as normal operation. To completely evaluate system sounds, a series of
templates should be used to test all normal sounds, and only after all templates are unable to match the sound would the system issue an error message.

   Optimizing histograms and modifying the code to allow multiple templates will be left as future work.

5.7 Summary

Since hard drives exhibit different failure modes, a high level of agreement between S.M.A.R.T. and SonMon was not expected. Optimizing the threshold parameter for SonMon produced better performance than expected in the restricted test being applied.

Another potential concern was a failure during the course of the experiment to optimize the histogram template for individual hard drives. Despite the fact that the hard drives were produced by different manufacturers and were of differing sizes and designs, the single-template classifications of good and bad hard drive conditions were acceptable: a customized template was not needed for each hard drive. This result was unexpected.

Only one type of sound was submitted to SonMon to evaluate the performance of the system; other kinds of sounds that SonMon may encounter in a computer were ignored. A production version of SonMon would need to recognize several normal modes of operation so that these conditions will not be erroneously identified as failures. More work is needed on SonMon in order for it to be applied as a practical tool to monitor computer systems.
CHAPTER 6

CONCLUSIONS

Based on the test results of SonMon, sound can be used to detect abnormal hardware operation. A production quality application to monitor sound remains an unrealized goal, but the prototype provided here demonstrates the viability of the monitoring approach.

The results also indicate that SonMon is not a stand-alone solution to the problem. Hard drive faults exist that are impossible for SonMon to detect. For example, electrical signal failures in which the drive runs and makes normal sounds but does not respond to the system will not be detected by SonMon. These types of failures are detected by a S.M.A.R.T. monitoring application such as smartmontools. It is recommended that S.M.A.R.T. monitoring tools such as smartmontools be used in conjunction with a sound monitoring hardware system such as SonMon.

Future Work

Error handling in SonMon should be improved. Currently, if the return code from 
\texttt{snd_pcm_readi} function indicates that an error has occurred, the application exits with a “broken pipe” error message. Typically, this error occurs with a sound card overrun where ALSA was not able to handle the incoming data before it was overwritten with newer data. SonMon should be modified to recover from overruns without interrupting the system. If an overrun occurs, the lost data are not critical to SonMon’s successful operation and the system can be restarted.

Currently, the histogram template that is compared to the wavelet output histogram is hard-coded. Hard coding the histograms template is bad practice and histogram data should be loaded at startup from a data file. The threshold value for chi-square should also be loaded from a file, as well as parameters such as, sampling rate, microphone volume control setting and other mixer settings, and IP or name of client machine where the error message should be sent.

Since the sound input level affects the accuracy of the detection system, the sound card’s microphone volume control should be controlled automatically to normalize input to the system.
The automatic volume control must be carefully designed so that the SonMon detection system would not mistake automatic volume control changes for hardware faults.

A “Leaky Bucket” algorithm should be considered for reporting failure. Instead of triggering on a single error, the system could allow a number of errors to accumulate over a period of time before signaling a failure.

It might also be beneficial to configure the production application as a daemon so that it can run in the background as a system service.

SonMon should be designed with standard Linux features such as command line options. SonMon should at minimum support the -h and --help options, which would display a help message. Other command line options allowing user configuration or adjustments may be added as may be desired by users. A command line option allowing multiple configuration files so that users could experiment with different histogram settings would be useful. A command line option to allow SonMon to export histogram data so that users could build their own histogram sets would be useful.

SonMon or its components should be investigated for suitability for use with other types of signals that may be present in a computer. Other signals present in a computer include electrical noise or sounds beyond the limits of the built-in sound card or microphone. The wavelet filtering class may be useful with these sounds.
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Sony Model C-22 FET Condenser Microphone Instruction Manual, Sony Corporation

*Speech Recognition using Daubechies Wavelets*


APPENDICES
Appendix A: Glossary

Filter—system of hardware or software that modifies the frequency content of a signal

Frequency-Domain—representation of signals as a function of frequency. Usually the result of a transformation function such as Fourier or Wavelet transform.

GNU—Project to develop free software for UNIX systems. The name GNU is a recursive acronym for “GNU’s Not UNIX.” More information on the GNU project can be found at http://www.gnu.org

GNU GPL—General Public License—Distribution license for most GNU software. Under the GPL, the software source code may be used freely, without royalty, and modified without permission. The GPL that software modified from a GPL source must also made available under the GPL. The text of the GPL may be seen at http://www.gnu.org/licenses/gpl.html

High-pass Filter—Filter that sends high frequency components of a signal to the output while attenuating low frequency components.

Low-pass Filter—Filter that sends low frequency components of a signal to the output while attenuating high frequency components.

Maple—A mathematics modeling program for solving mathematical problems and creating interactive technical applications. A product of Waterloo Maple, Inc.

http://www.maplesoft.com

Quadrature Mirror Filter (QMF)—Symmetric filter bank which splits signals into two bands. Wavelet analysis via multiresolution analysis constitutes such a filter.

Rib—Individual datum on a wavelet ridge

Ridge—Curve formed in the wavelet time-frequency plane by maxima in the wavelet transform

Template—In SonMon, a set of histogram values being used for comparison with the wavelet data being read. Not to be confused with a C++ template.

Time-domain—representation of signals as a function of time. The usual way of representing signals. Sometimes referred to as spatial-domain.
Appendix B: Source Code

D4.h

/* declaration file for D4 wavelet multi-resolution analysis

Programmer: Robert K. Henry
Project: D4
Date created: 11/25/2004
Last modified: 8/18/2005
Purpose: An implementation of the Daubechies D4 wavelet multi-resolution analysis

Description:
Data are added into the member item "data" using the method "add()". Since the size of "data" is 1024, no more 1024 elements can be added.

When data array is full with 1024 items, indicated by the counter "index," then the transformation function "transform()" may be invoked. Transform will not function unless 1024 elements have been loaded.

The method "transform()" performs a forward D4 transform using multi-resolution analysis. Since there are 1024 elements in the input, there will be 10 levels of analysis. No inverse transform is provided.

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*/

class d4
{
    public:
        double getMax(int level); //Returns the largest element in specified transform level
        double getMin(int level); //Returns the smallest element in specified transform level
        double getLevelSq(int item, int level); //Squares values before returning them
double getLevel(int item, int level); // accessor for data by level
double getRaw(int a); // Accessor for data in unsorted order
int maxIndex(int level); // Returns number of elements for level
void reset(); // initializes for reloading with new data
int transform(); // performs forward transform
void add(double item); // adds one data item to data
d4();
virtual ~d4();

protected:
const double p0, p1, p2, p3;
unsigned int index; // number of elements currently in data array
double data[1024]; // input data and storage for transform
double maxValue[10]; // maximum value in specified level
double minValue[10]; // minimum value in specified level
bool already_transformed;
/* Implementation of D4 Multi-resolution analysis

Programmer: Robert K. Henry
Project: D4
Date created: 11/25/2004
Last modified: 8/18/2005
Purpose: An implementation of the Daubechies D4 wavelet multi-resolution analysis

Description:
Data are added into the member item "data" using the method "add()". Since the size of "data" is 1024, no more 1024 elements can be added.

When data is full with 1024 items, indicated by the counter "index," then the transformation function "transform()" may be invoked. Transform will not function unless 1024 elements have been loaded.

The method "transform()" performs a forward D4 transform using multi-resolution analysis. Since there are 1024 elements in the input, there will be 10 levels of analysis.

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*/

#include "D4.h"

d4::d4():p0(0.683012702), p1(1.183012702), p2(0.316987298), p3(-0.183012702),
index(0), already_transformed(false)
{
}

d4::~d4()
{
}
// Function name : d4::add
// Description     : Adds one data item to data array.
// Return type  : void
// Argument         : double item
void d4::add(double item)
{
    if (index < 1024)
    {
        data[index] = item;
        index++;
    }
}

// Function name : d4::transform
// Description     : Performs a forward wavelet transform using the Daubechies
//                    D4 and multiresolution analysis
//                    Data is left in the original array
//                    Original data is destroyed.
//                    1024 elements must have already been loaded using add()
// Return type  : int,
//                0 if successful (1024 elements were available)
//                -1 if unsuccessful (input wasn't full with 1024 elements)
//                or transform has already been executed on data
int d4::transform()
{
    double a0,a1,a2,a3;
    if ((index == 1024) && (!already_transformed))
    {
        unsigned int step = 1;
        unsigned int increment = 2;
        unsigned int level = 0; //used as index for min & max
        while (step < index)
        {
            for ( i=0;i<1024;i+=increment)
            {
                //load temp data
                a0 = data[i];
                a1 = data[i + step];

                // For last elements, end of data may reach beyond end of array
                // if so, enter zeros for these two elements
                if ((i + 2 * step) > 1023)
                {
                    a2 = 0;
                }
                else
                {
                    a2 = data[i + 2 * step];
                }
            }
        }
    }
}
// Last two elements are checked independently
if ((i + 3 * step) > 1023)
{
    a3 = 0;
} else
{
    a3 = data[i + 3 * step];
}

//low-pass filter
data[i] = (a0 * p0 +
    a1 * p1 +
    a2 * p2 +
    a3 * p3) / 2;

//high-pass filter--detail
data[i+step] = (a0 * p3 -
    a1 * p2 +
    a2 * p1 -
    a3 * p0) / 2;

//find min & max
if (i > 0) //if this is the first number just copy the first number
{
    if (maxValue[level] < data[i+step]) maxValue[level] = data[i+step];
    if (minValue[level] > data[i+step]) minValue[level] = data[i+step];
    else
    {
        minValue[level] = maxValue[level] = data[i+step];
    }
}
step *= 2;
increment *= 2;
level ++;
already_transformed = true;
return 0;
}
else
{
    return -1;
}
// Function name: d4::reset
// Description: Resets the state of the object to initial state with
//      no elements loaded.
// Return type: void
void d4::reset() {
  index = 0;
  already_transformed = false;
}

// Function name: d4::getRandom
// Description: Returns data item at random from the specified level
//      in the tree.
//       Check the number of items at the specified level using
//      maxIndex().
// Return type: double
// Argument: int item
// Argument: int level
double d4::getLevel(int item, int level) {
  int i = 1;
  int b = i << level;
  int a = b << 1;
  if (item*a+b > 1024) throw ("D4: Index out of range");
  return data[item*a+b];
}

double d4::getLevelSq(int item, int level) {
  int i = 1;
  int b = i << level;
  int a = b << 1;
  if (item*a+b > 1024) throw ("D4: Index out of range");
  return data[item*a+b] * data[item*a+b];
}

double d4::getRaw(int a) {
  return data[a];
}
// Function name : d4::maxIndex
// Description : Returns the maximum index number of elements at
// each level of the tree
// Return type : int -- maximum number of elements at this level
// Argument : int level
int d4::maxIndex(int level)
{
    int i = 1024;
    return i >> (level + 1);
}

double d4::getMax(int level)
{
    if (level > 10) level = 10;
    if (level < 0) level = 0;
    return maxValue[level];
}

double d4::getMin(int level)
{
    if (level > 10) level = 10;
    if (level < 0) level = 0;
    return minValue[level];
}
sonmon.c

Main Program

/* Read data from sound card using ALSA under Linux
   Programmer: Robert K. Henry
   Project: SonMon
   Date created: 5/11/2005
   Last modified: 8/18/2005
   Purpose: Read data from sound card in Linux using ALSA
           Asynchronous technique using callback function
   Description:
       Main opens a PCM device under ALSA and sets the hardware
       and software configurations.

       Main then registers a callback function,
       async_callback().

       When the sound card capture device has received
       1024 data elements, or "frames"
       the callback function activates and extracts the
       captured data and stores
       it in the array buf. Buf consists of 16-bit signed
       numbers so the sound card is
       configured to deliver 16-bit little-endian values using the
       parameter SND_PCM_FORMAT_S16_LE.

       When the callback function has extracted 1024 frames,
       the frames are then
       processed.

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   http://equalarea.com/paul/alsa-audio.html#interruptex
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*/
```c
#include <stdio.h>
#include <unistd.h>
#include <stdlib.h>
#include <errno.h>
#include <alsa/asoundlib.h>
#include "D4.h"
#include <iostream>
#include <fstream>
#include <string>
#include <string.h>
using namespace std;

snd_pcm_t *capture_handle;
short buf[4096];
char debugstring[30] = "No Frames Read";
char success[30] = "Frames successfully read";
char *device = "plughw:0,0"; /* sound device */
unsigned int rate = 44100; /* stream rate */
int count = 0; // used for debugging
int icounter; // used for debugging
int lcounter; // and index for counting wavelet levels
int i, j, k;
int score;
int scoreThreshold = 1000;
int displayState = 0;

d4 w; // a wavelet object
snd_pcm_uframes_t buffer_size;
snd_pcm_uframes_t period_size = 1024;
double bins[12] = {64, 128, 512, 1024, 2048, 4096, 8192, 16384,
                   32768, 65536, 131072};
int histogram[13];
int histogramTemplate[6][13] = {
    {135, 57, 76, 83, 141, 20, 0, 0, 0, 0, 0, 0, 0},
    {32, 8, 13, 22, 70, 45, 40, 17, 9, 0, 0, 0, 0},
    {12, 6, 6, 7, 26, 19, 26, 19, 6, 1, 0, 0, 0},
    {13, 6, 6, 11, 19, 8, 1, 0, 0, 0, 0, 0, 0},
    {7, 4, 6, 4, 9, 2, 0, 0, 0, 0, 0, 0, 0},
    {5, 3, 2, 3, 3, 0, 0, 0, 0, 0, 0, 0, 0}};

// fstream fout; // uncomment to export files

struct async_private_data {
    signed short *samples;
    snd_pcm_channel_area_t *areas;
    double phase;
};

int classify ()
{
    double element;
    int chi;
    int count;
    chi = 0;
```
for (j=0; j<6; j++)
{
  //clear histogram to accept a new count for this level of the wavelet
  for (i = 0; i < 13; i++) histogram[i] = 0;
  //step through this level of the wavelet
  for (i=0; i<w.maxIndex(j); i++)
  {
    //get an element from the wavelet and store it locally
    element = w.getLevelSq(i,j);
    k = 0;

    //find which bin the element fits into
    while ((element > bins[k]) && (k < 12))
    {
      k++;
    }
    //increment the count for the histogram bin which matched--or the
    last bin if it didn't match
    histogram[k]++;
  }

  // print histogram values to file
  //***Uncomment the next 5 lines to export histogram to a file.
  //fout << "level = " << j << endl;
  //for (i=0;i<13;i++)
  //{ // fout << histogram[i] << ", ";
  //}
  //fout << "end" << endl;

  // go through the histogram and calculate chi-square scores
  // for each bin, adding it to the total
  for (i=0;i<13;i++)
  {
    // if the template value is zero, skip this part to avoid
    // a divide by zero error
    if (histogramTemplate[j][i] > 0)
    {
      count = histogram[i]-histogramTemplate[j][i];
      chi += count*count/histogramTemplate[j][i];
    }
  }
  //fout << chi << endl; //uncomment to export chi-square score to a file
  return chi; //return the chi-square total
}

// Function name : async_callback
// Description :
// Return type : void
// Argument : ahandler, a pointer to the function's arguments
// Description : An ALSA callback. When ALSA is interrupted by
// the sound card,
// ALSA executes this function to handle the capture data.
static void async_callback(snd_async_handler_t *ahandler) {
  snd_pcm_t *handle = snd_async_handler_get_pcm(ahandler);
// short *data = snd_async_handler_get_callback_private(ahandler);
// signed short *samples = data->samples;
snd_pcm_sframes_t avail;
int err;
avail = snd_pcm_avail_update(handle);
while (avail >= period_size) {
   .err = snd_pcm_readi(handle, buf, period_size);
    if (err < 0) {
        printf("Read error: %s\n", snd_strerror(err));
        exit(EXIT_FAILURE);
    }
    if (err != period_size) {
        printf("Read error: Read %li expected %li\n", err, period_size);
        exit(EXIT_FAILURE);
    }
    w.reset();
    for (icounter = 0; icounter < 1024; icounter++) w.add(buf[icounter]);
    w.transform();
    //for (icounter = 0; icounter< w.maxIndex(4);icounter++)
    //  printf(" %4.2f",w.getLevel(icounter,4));
    score = classify();
    //printf("%d\n", score);
    if (score > scoreThreshold) {
        if (displayState == 0) {
            printf("Chi-square exceeds %d\n", score);
            displayState = 1;
        }
    } else {
        if (displayState == 1) {
            printf("Chi-square ok %d\n", score);
            displayState = 0;
        }
    }
    //printf("end\n");
    avail = snd_pcm_avail_update(handle);
}

main(int argc, char *argv[]) {
    //fprintf(stderr, "starting program\n");
    snd_pcm_hw_params_t *hw_params;
    snd_pcm_sw_params_t *sw_params;
    snd_pcm_sframes_t frames_to_deliv;
    int nfds;
    int err;
Instead of using SND_PCM_NONBLOCK or SND_PCM_ASYNC for the last value
of snd_pcm_open, the last value must be "0".
Otherwise the card won't start.

```c
if ((err = snd_pcm_open(&capture_handle, device,
SND_PCM_STREAM_CAPTURE,0)) < 0) {
    fprintf(stderr, "cannot open audio device %s (%s)\n",
    device,
    snd_strerror(err));
    exit(1);
}

//Allocate a hw_params element
if ((err = snd_pcm_hw_params_malloc(&hw_params)) < 0) {
    fprintf(stderr, "cannot allocate hardware parameter structure
%s\n",
    snd_strerror(err));
    exit(1);
}

//Initialize the hw_params element with default values.
if ((err = snd_pcm_hw_params_any(capture_handle, hw_params)) < 0) {
    fprintf(stderr, "cannot initialize hardware parameter structure
%s\n",
    snd_strerror(err));
    exit(1);
}

//Set access type to interleaved. There will only be 1 channel but
interleaved will work anyway
if ((err = snd_pcm_hw_params_set_access(capture_handle, hw_params,
SND_PCM_ACCESS_RW_INTERLEAVED)) < 0) {
    fprintf(stderr, "cannot set access type (%s)\n",
    snd_strerror(err));
    exit(1);
}

//Set format to 16-bit signed little-endian. Little endian will work on
//an Intel machine but this might need to be changed for other hardware.
//16-bit signed is expected for the DSP system. A larger bit size might
//work on some hardware but might not be portable
if ((err = snd_pcm_hw_params_set_format(capture_handle, hw_params,
SND_PCM_FORMAT_S16_LE)) < 0) {
    fprintf(stderr, "cannot set sample format (%s)\n",
    snd_strerror(err));
    exit(1);
}

//set the rate. Rate will be user-configurable
if ((err = snd_pcm_hw_params_set_rate_near(capture_handle, hw_params,
&rate, 0)) < 0) {
    fprintf(stderr, "cannot set sample rate (%s)\n",
    snd_strerror(err));
    exit(1);
}
```

// Set the number of channels to 1. Only one channel will be analyzed.
if ((err = snd_pcm_hw_params_set_channels(capture_handle, hw_params, 1)) < 0) {
    fprintf(stderr, "cannot set channel count (%s)\n",
            snd_strerror(err));
    exit(1);
}

// Set the above specified parameters
if ((err = snd_pcm_hw_params(capture_handle, hw_params)) < 0) {
    fprintf(stderr, "cannot set parameters (%s)\n",
            snd_strerror(err));
    exit(1);
}

snd_pcm_hw_params_free(hw_params);

    /* tell ALSA to wake us up whenever 1024 or more frames
     of playback data can be delivered. Also, tell
     ALSA that we'll start the device ourselves.
    */

    // Allocate space for a sw_params data element
if ((err = snd_pcm_sw_params_malloc(&sw_params)) < 0) {
    fprintf(stderr, "cannot allocate software parameters structure
            (%s)\n",
            snd_strerror(err));
    exit(1);
}

    // snd_pcm_sw_params_current initializes sw_params with current software
    // settings
if ((err = snd_pcm_sw_params_current(capture_handle, sw_params)) < 0) {
    fprintf(stderr, "cannot initialize software parameters structure
            (%s)\n",
            snd_strerror(err));
    exit(1);
}

    // Set avail min inside a software configuration container.
    // Sets Minimum avail frames to consider PCM ready, 1024
if ((err = snd_pcm_sw_params_set_avail_min(capture_handle, sw_params,
                                          1024)) < 0) {
    fprintf(stderr, "cannot set minimum available count (%s)\n",
            snd_strerror(err));
    exit(1);
}

    // Sets start threshold in frames. 0 in this case.
    // PCM is automatically started when playback frames available to PCM are
    // >= threshold
    // or when requested capture frames are >= threshold
if ((err = snd_pcm_sw_params_set_start_threshold(capture_handle,
                                                sw_params, 0)) < 0) {
    fprintf(stderr, "cannot set start mode (%s)\n",
            snd_strerror(err));
    exit(1);
}
//Install PCM software configuration defined by params.
if ((err = snd_pcm_sw_params(capture_handle, sw_params)) < 0) {
    fprintf(stderr, "cannot set software parameters (%s)\n",
            snd_strerror(err));
    exit(1);
}

/* the interface will interrupt the kernel every 1024 frames, and ALSA
will wake up this program very soon after that. */

//Prepare PCM for use.
if ((err = snd_pcm_prepare(capture_handle)) < 0) {
    fprintf(stderr, "cannot prepare audio interface for use (%s)\n",
            snd_strerror(err));
    exit(1);
}

//struct async_private_data data;
snd_async_handler_t *ahandler;
int count;

period_size = 1024;
//printf("setting up callback\n");
err = snd_async_add_pcm_handler(&ahandler, capture_handle,
    async_callback, &buf);
    if (err < 0) {
        printf("Unable to register async handler\n");
        exit(EXIT_FAILURE);
    }

err = snd_pcm_start(capture_handle);
if (err < 0) {
    printf("Start error: %s\n", snd_strerror(err));
    exit(EXIT_FAILURE);
}

//open output file
//fout.open("D4Output.txt", ios::out); //uncomment to export data to file

//display copyright notice
printf("SonMon, Copyright (C) 2005 Robert Henry\nSonMon comes with
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/* because all other work is done in the signal handler,
   suspend the process */
while (1) {
    sleep(1);
}
printf(debugstring);
snd_pcm_close(capture_handle);
//fout.close();       //uncomment if exporting to a file.
exit(0);
## Makefile

A simple make file for use in Linux to build the system

```
all: sonmon

sonmon: sonmon.o D4.o
    g++ sonmon.o D4.o -lasound -o sonmon

D4.o: D4.cpp
    g++ D4.cpp -c

sonmon.o: sonmon.c
    g++ sonmon.c -c
```
Builds the wavelet into a DLL suitable for use with the mathematics analysis application Maple.

/**
  Programmer: Robert K. Henry
  Project: D4
  Date created: 12/4/2004
  Last modified: 8/2/2005
  Purpose: Interface the D4 wavelet to DLL for Maple
  Description:
  Data are supplied in the array input. Input is fixed at 1024 elements. The input array MUST be defined in Maple as a c_order array in row-major order. Maple arrays default to Fortran_order or column-major order. In a 1-dimensional array the two configurations will be the same. The output array MUST be declared in Maple using the Maple C_order option.

  The DLL copies the data to the wavelet class and invokes the transform function. The resulting transform is returned in the two-dimensional array "output". Like the input array, this array must also be a C-style row-major order array and is pre-defined in Maple. The array must be 512x10 but only half these elements will be used since the output array is triangular.

  Compiling:
  to compile with Microsoft VC++, first set the environment variables with vcvars32.bat

  Then invoke the compiler with
  cl d4dll.cpp D4.cpp -Gz -LD -link -export:d44Maple

  Using in Maple:
  The DLL must be declared in Maple. This can be done with:
  > d4_CTrans :=
  > define_external("d44Maple", LIB="d4trans.dll",
  > input::ARRAY(1..1024,float[8]),
  > output::ARRAY(1..512,1..10,float[8])):

  where the name of the DLL is "d4trans.dll"

  The output array MUST be declared as C_order like this:
  d1C:=Matrix(1..10,1..512,datatype=float[8],order=C_order):
  */
#include "D4.h"
declspec(dllexport) void d44Maple( double input[1024], double output[10][512])
{
  // it is assumed that the input array is 1024
  // and the output array is 10 x 512
  int i, level;
  double t;
d4 h;
    h.reset(); // to make sure the object starts fresh each time
    for (i = 0; i < 1024; i++) h.add(input[i]);
    h.transform();

    for (level = 0; level < 10; level++){
        for (i = 0; i < h.maxIndex(level); i++){
            t = h.getLevelSq(i, level);
            output[level+i*10] = t;
            output[level][i] = t;
        }
    }
Appendix C: Gnu General Public License

Version 2, June 1991

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END OF TERMS AND CONDITIONS
### Appendix D: Results Summary

<table>
<thead>
<tr>
<th>Manufacturer</th>
<th>Drive ID</th>
<th>S.M.A.R.T. Results</th>
<th>Total Frames</th>
<th>Error Count</th>
<th>%</th>
<th>Error Count</th>
<th>%</th>
</tr>
</thead>
<tbody>
<tr>
<td>Connor</td>
<td>CA905DG</td>
<td>No Response</td>
<td>969</td>
<td>2</td>
<td>0.21</td>
<td>13</td>
<td>1.34</td>
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<tr>
<td>IBM</td>
<td>TVLF1297</td>
<td>No Response</td>
<td>883</td>
<td>1</td>
<td>0.11</td>
<td>18</td>
<td>2.04</td>
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<tr>
<td>Maxtor</td>
<td>00990202-19661-76P-H2UH</td>
<td>No Response</td>
<td>1047</td>
<td>2</td>
<td>0.19</td>
<td>2</td>
<td>0.19</td>
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<tr>
<td>Maxtor</td>
<td>A348F6VCDA43A</td>
<td>No S.M.A.R.T.</td>
<td>1002</td>
<td>154</td>
<td>15.37</td>
<td>264</td>
<td>26.35</td>
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<tr>
<td>Maxtor</td>
<td>E40AR3LA</td>
<td>FAILED Seek_Error_Rate</td>
<td>744</td>
<td>20</td>
<td>2.69</td>
<td>99</td>
<td>13.31</td>
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<td>Maxtor</td>
<td>E40BFG9A</td>
<td>No Response</td>
<td>981</td>
<td>981</td>
<td>100</td>
<td>981</td>
<td>100</td>
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<tr>
<td>Maxtor</td>
<td>K309T6JA</td>
<td>Passed</td>
<td>905</td>
<td>0</td>
<td>-</td>
<td>0</td>
<td>-</td>
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<tr>
<td>Maxtor</td>
<td>K607NJ1A</td>
<td>Passed</td>
<td>793</td>
<td>0</td>
<td>-</td>
<td>0</td>
<td>-</td>
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<tr>
<td>Maxtor</td>
<td>L306XGDA</td>
<td>No S.M.A.R.T.</td>
<td>967</td>
<td>18</td>
<td>1.86</td>
<td>48</td>
<td>4.96</td>
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<tr>
<td>Maxtor</td>
<td>T1RNZVFC</td>
<td>No Response</td>
<td>948</td>
<td>0</td>
<td>-</td>
<td>1</td>
<td>0.11</td>
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<tr>
<td>Maxtor</td>
<td>W40S522VA</td>
<td>Passed</td>
<td>837</td>
<td>0</td>
<td>-</td>
<td>0</td>
<td>-</td>
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<tr>
<td>Quantum</td>
<td>79112173912</td>
<td>No Response</td>
<td>1008</td>
<td>0</td>
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<td>97X1802B</td>
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<td>Quantum</td>
<td>SG-035NTD-12547-15E-15RG</td>
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<td>0.12</td>
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<td>Western Digital</td>
<td>WM 342 197 3562</td>
<td>No Response</td>
<td>1072</td>
<td>382</td>
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<td>530</td>
<td>49.44</td>
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<td>Western Digital</td>
<td>WM 356 043 2191</td>
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<td>977</td>
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<td>0.41</td>
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<td>WM 361 044 2859</td>
<td>FAILED Raw_Read_Error_Rate</td>
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<td>0</td>
<td>-</td>
<td>3</td>
<td>0.45</td>
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<td>Western Digital</td>
<td>WM 361 214 3523</td>
<td>No S.M.A.R.T.</td>
<td>711</td>
<td>0</td>
<td>-</td>
<td>1</td>
<td>0.14</td>
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<td>Western Digital</td>
<td>WM 408 031 9718</td>
<td>Passed</td>
<td>1318</td>
<td>1</td>
<td>0.08</td>
<td>2</td>
<td>0.15</td>
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<td>Western Digital</td>
<td>WM 408 032 8026</td>
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<td>583</td>
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<td>50.94</td>
<td>351</td>
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<td>Western Digital</td>
<td>WM 408 071 0209</td>
<td>Passed—history of Raw_Read_Error_Rate</td>
<td>674</td>
<td>2</td>
<td>0.30</td>
<td>17</td>
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<td>Western Digital</td>
<td>WM 627 362 1072</td>
<td>Passed</td>
<td>879</td>
<td>5</td>
<td>0.57</td>
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<td>2.84</td>
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<td>Western Digital</td>
<td>WM 627 362 7493</td>
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<td>922</td>
<td>0</td>
<td>-</td>
<td>3</td>
<td>0.33</td>
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<td>Western Digital</td>
<td>WS 342 200 0122</td>
<td>No S.M.A.R.T.</td>
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<td>23.95</td>
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<td>Western Digital</td>
<td>WT 336 017 2647</td>
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<td>659</td>
<td>0</td>
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<td>Western Digital</td>
<td>WT 359 004 9279</td>
<td>No Response</td>
<td>858</td>
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<td>75</td>
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<td>Manufacturer</td>
<td>Drive ID</td>
<td>S.M.A.R.T. Results</td>
<td>Total Frames</td>
<td>Error Count</td>
<td>Error Count</td>
<td>Error Count</td>
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<td>Western Digital</td>
<td>WT 473 311 2127</td>
<td>No S.M.A.R.T.</td>
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<td>Western Digital</td>
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<td>942</td>
<td>0</td>
<td>-</td>
<td>0</td>
<td>-</td>
</tr>
</tbody>
</table>
VITA
ROBERT K. HENRY

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Marital Status: Single

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